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Line-graph qubit routing: from kagome to heavy-hex and more

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Quantum computers have the potential to outperform classical computers, but are currently limited in their capabilities. One such limitation is the restricted connectivity between qubits, as captured by the hardware's coupling graph. This limitation poses a challenge for running algorithms that require a coupling graph different from what the hardware can provide. To overcome this challenge and fully utilize the hardware, efficient qubit routing strategies are necessary. In this paper, we introduce line-graph qubit routing, a general method for routing qubits when the algorithm's coupling graph is a line graph and the hardware coupling graph is a heavy graph. Line-graph qubit routing is fast, deterministic, and effective; it requires a classical computational cost that scales at most quadratically with the number of gates in the original circuit, while producing a circuit with a SWAP overhead of at most two times the number of two-qubit gates in the original circuit. We implement line-graph qubit routing and demonstrate its effectiveness in mapping quantum circuits on kagome, checkerboard, and shuriken lattices to hardware with heavy-hex, heavy-square, and heavy-square-octagon coupling graphs, respectively. Benchmarking shows the ability of line-graph qubit routing to outperform established general-purpose methods, both in the required classical wall-clock time and in the quality of the solution that is found. Line-graph qubit routing has direct applications in the quantum simulation of lattice-based models and aids the exploration of the capabilities of near-term quantum hardware.

I. INTRODUCTION

Quantum computing offers potential to revolutionize a wide range of domains by efficiently solving problems that are intractable for classical computers [1, 2]. To run any quantum algorithm, it must be compiled into a quantum circuit that can be executed on the quantum hardware. The hardware coupling graph of a quantum computer, which defines adjacency between qubits based on the ability to perform two-qubit gates between them, plays a crucial role in this process. The problem of ensuring that a quantum circuit is compatible with the hardware coupling graph is referred to as the qubit routing problem [3, 4]. While generalized methods exist for qubit routing [5], a standard approach to implement two-qubit gates between non-adjacent qubits is to insert SWAP gates, making the qubits effectively adjacent [3, 4, 6-9]. To obtain a practical quantum advantage on noisy intermediate-scale quantum (NISQ) [10] devices, it is imperative that overhead arising from compilation is kept to a minimum [11, 12]. However, finding the swapping strategy that requires the least number of SWAP gates is NP-hard [3, 6], making it a challenging problem to solve. Heuristic, probabilistic methods have been developed, but their classical runtime may become problematic for large circuits, and the solution they find may be far from optimal [4, 7–9]. Striking the right balance between the classical resources required for routing and minimizing the circuit depth of the routed quantum circuit is crucial in maximizing the performance of NISO devices.

The qubit routing problem is particularly evident in the quantum simulation of lattice-based spin models. One of the first areas in which it was realized that quantum computers could outperform classical computers was that of quantum simulation [13]. When applied specifically to lattice-based spin systems with two-body interactions, quantum simulation by Trotterization [14] approximates the overall time evolution operator of the quantum-mechanical system by a sequence of two-qubit gates, where each two-qubit gate corresponds to the time evolution according to one two-body term in the Hamiltonian (dynamic quantum simulation) [15, 16]. Additionally, by introducing variable parameters for the per-term evolution times, these circuits are transformed to circuits that prepare ansatz states for the variational quantum eigensolver (VQE) [17], designed to variationally find the ground state of the quantum-mechanical system (static quantum simulation). Before any routing, these circuits for dynamic and static quantum simulation naturally require hardware with a coupling graph that is equal to the lattice of the lattice-based spin system (the virtual graph) [17]. There will generally be a mismatch between the virtual graph and the hardware coupling graph. Efficient qubit routing plays a crucial role in overcoming this mismatch.

In this paper, we develop an efficient and deterministic qubit routing strategy, which we call line-graph qubit routing, or line-graph routing for short. It maps any circuit on a line graph L(G) to hardware with coupling graph heavy(G). Here, heavy(G) is obtained from the graph G by placing a node on every edge of G. We call these added nodes the heavy nodes of heavy(G). By definition, the nodes of the line graph L(G) consist of the heavy nodes of heavy(G). In L(G), two nodes are adjacent if the associated edges in G are incident on the same



Table I. Examples of line-graph routing. Line-graph routing maps any circuit with coupling graph L(G) (blue edges) to circuits with coupling graph heavy(G) (black edges). Line-graph routing finds direct application in the quantum simulation of the magnetic properties of some (model) materials with coupling graph L(G) (last row) on hardware with coupling graph heavy(G).

node of G. It is instructive to verify this property for one of the pairs (L(G), heavy(G)) in Table I. For example, line-graph routing maps the circuits for the quantum simulation of the Heisenberg anti-ferromagnet (HAFM) on the kagome lattice to hardware with a heavy-hex coupling graph. In this example, G is the hexagonal lattice and L(G) is the kagome lattice. Despite these examples, we stress that line-graph routing is applicable to any circuit on any line graph L(G).

The remainder of this paper is organized as follows. We first introduce line-graph routing by example, mapping circuits for the quantum simulation of the kagome HAFM to hardware with a heavy-hex coupling graph (Sec. II). We formalize and generalize this approach to arbitrary circuits and arbitrary line graphs in Sec. III. We benchmark our software implementation of the general algorithm against existing qubit routing approaches in Sec. IV, to conclude with a discussion and outlook in Sec. V.

II. KAGOME TO HEAVY-HEX

Line-graph routing is arguably most clearly explained with an example, which we do in this section by mapping circuits for the quantum simulation of the kagome HAFM to quantum hardware with a heavy-hex coupling graph. First, we use this example due to the relevance of the kagome spin model in exploring quantum phenomena like topological states of matter and quantum spin liquids [21, 27]. The kagome lattice's significance extends to chemistry, as it is frequently observed in transition metal compounds and metal organic frameworks [28, 29]. The ground state of the kagome HAFM is a long-standing open problem in quantum magnetism [30] that can potentially be solved on NISQ devices [16]. By classical emulation of noiseless quantum computers, it was previously demonstrated that the ground-state energy found by a VQE approaches the true ground-state energy exponentially as a function of the circuit depth [16].

Second, the heavy-hex coupling graph is the coupling graph of IBM's current and future superconducting hardware [31, 32]. Among the emerging quantum hardware platforms, IBM's superconducting qubits have gained significant attention due to their rapid development and scalability in the NISQ era. To optimize this superconducting qubit hardware and mitigate the occurrence of frequency collisions and crosstalk [33–35], error correcting codes are designed on low-degree graphs such as heavy-hex and heavy-square lattices, preventing errors during program execution [31, 32, 36]. This motivates the further development of hardware with these types of connectivity graphs.

The relevance of the kagome-to-heavy-hex mapping was further highlighted by the IBM Quantum's Open Science Prize 2022, where the challenge was to prepare the ground state of the kagome HAFM using a VQE and implement it on a 16-qubit IBM Quantum Falcon device with a heavy-hex coupling graph [37]. It is important to note that, also within the context of quantum simulation, the routing problem is not unique to the quantum simulation of spin problems on the kagome lattice. Other lattice-based spin models, such the HAFM on the shuriken lattice, are also known for their geometric frustration and challenging simulation [38].

A. Line-graph routing

The first step in line-graph routing is establishing a one-to-one correspondence between the nodes of the virtual graph (in this section, the kagome lattice) and the hardware connectivity graph (in this section, the heavyhex lattice). This correspondence is achieved by aligning the nodes of the kagome lattice with the heavy nodes of the heavy-hex lattice, as shown in Fig. 1. Subsequently, the light (non-heavy) nodes of the heavy-hex lattice are used to mediate two-qubit gates between the spins on the nodes of the kagome lattice.

To see this in more detail, assume a kagome quantum circuit, that is, a circuit composed of single-qubit gates on qubits $\{i\}$ and two-qubit gates along the edges $\{(i, j)\}$ of a patch of the kagome lattice. To map the circuit from the kagome to the heavy-hex lattice, we label the heavy qubits on the heavy-hex lattice with the labels $\{i\}$ of the congruent qubits on the kagome lattice, as shown in Fig. 1. Under this identification, any single-qubit gate in the kagome circuit is trivially mapped to a single-qubit gate on the heavy-hex lattice.

To map the two-qubit gates, let us label the ℓ th twoqubit gate in the kagome circuit, acting on qubits (i, j), by U_{ij}^{ℓ} . Any such gate can be performed on the heavyhex lattice by mapping it to the three-qubit *mediated two-qubit gate MU*

$$U_{ij}^{\ell} \mapsto M U_{imj}^{\ell} = \mathrm{SWAP}_{mi} U_{mj}^{\ell} \mathrm{SWAP}_{im}, \qquad (1)$$

where qubit $m = m_{ij}$ mediates the interaction between qubits *i* and *j*. This map provides the cornerstone of line-graph routing. The key point of line-graph routing is that for every pair of qubits (i, j) the existence and uniqueness of the mediating qubit $m = m_{ij}$ is guaranteed by the definition of L(G) and heavy(G) (see Sec. I).

Equation (1) introduces many SWAP gates that need not be performed physically. First, SWAP gates occurring at the beginning and end of the routed circuit can be accounted for by a relabeling of the qubits. Second, any two consecutive SWAP gates can be cancelled. These double SWAP gates are introduced by Eq. (1) if there are two consecutive two-qubit gates acting on the same two qubits (possibly with additional single-qubit gates on those qubits in between). Double SWAP gates are also introduced by Eq. (1) in the case of two consecutive two-qubit gates that have a single qubit in common and where the two resulting mediated gates have a mediating qubit in common. That is, if $m = m_{ij} = m_{ik}$, we have by Eq. (1) that

$$U_{ik}^{\ell'}U_{ij}^{\ell} \mapsto (\text{SWAP}_{mi}U_{mk}^{\ell'}\text{SWAP}_{im})(\text{SWAP}_{mi}U_{mj}^{\ell}\text{SWAP}_{im}) = \text{SWAP}_{mi}U_{mk}^{\ell'}U_{mj}^{\ell}\text{SWAP}_{im}.$$
(2)

In summary, line-graph routing first associates the nodes of L(G) with the heavy nodes of heavy(G), applies the map of Eq. (1) to all two-qubit gates, and finally removes superfluous SWAP gates as described above.

B. Application: quantum simulation

One immediate application of the kagome-to-heavyhex mapping is in the quantum simulation of the kagome



Figure 1. The kagome lattice (colored edges) is the line graph of the hexagonal lattice (black edges). As shown in the figure, the nodes of the kagome lattice can be identified with the heavy nodes of the heavy-hex lattice. The colored edges represent one possible circuit with a kagome coupling graph. Every color represents a layer in this circuit. In the first layer, singlet states are created along the blue lines. Thereafter, HEIS gates [Eq. (5)] are applied along all colored edges in sequence, defining one circuit cycle. This cycle is repeated to obtain the complete circuit. Figure adapted from Ref. [16].

HAFM on heavy-hex quantum hardware. In units where $\hbar = 1$, the HAFM has a Hamiltonian

$$H = \sum_{(i,j)} H_{ij}, \qquad H_{ij} = X_i X_j + Y_i Y_j + Z_i Z_j, \qquad (3)$$

where the sum is over all edges (i, j) of a given graph. In the current section, this graph could be any patch of the kagome lattice. Here, X_i denotes the Pauli-X operator acting on qubit *i* (similarly for Y_i, Z_i). This Hamiltonian is straightforwardly generalized to arbitrary two-spin interactions along the edges of a graph [39], which would conceptually not change the constructions that follows.

The goal of dynamic quantum simulation is to compute expectation values of observables with respect to the time-evolved state $|\psi(t)\rangle = e^{-iHt} |\psi(0)\rangle$, given some initial state $|\psi(0)\rangle$. On a quantum computer, this can be achieved by applying the unitary $\prod_{ij} e^{-iH_{ij}(t/r)}$ to $|\psi(0)\rangle$ (the latter of which is assumed to be easy to prepare) a total of r times. That is,

$$|\psi(t)\rangle \approx \left(\prod_{ij} e^{-i\frac{t}{r}H_{ij}}\right)^r |\psi(0)\rangle.$$
 (4)

The error in this approximation is of the order t^2/r [14], assuming a perfect quantum computer. After the preparation of $|\psi(t)\rangle$, expectation values can be extracted by repeated preparation and measurements. Note that H_{ij} acts on two qubits, so that $e^{-iH_{ij}(t/r)}$ is a two-qubit unitary that can be decomposed into a few one and twoqubit gates that act on qubits i, j only. In the case of the HAFM, $e^{-iH_{ij}(t/r)}$ is called the HEIS gate [16]

$$\text{HEIS}_{ij}(\alpha) \equiv e^{-i\alpha/4} e^{-i\alpha H_{ij}/4},\tag{5}$$



Figure 2. The line-graph routed quantum circuit for the quantum simulation of the kagome HAFM (Fig. 1), focusing on one triangle. The circuits involving the other triangles are similar. Empty colored boxes indicate subcircuits coming in from other triangles. The white SWAP gates can be omitted by a relabeling of the qubits (first SWAP gate) and the cancellation of double SWAP gates [other white SWAP gates, cf. Eq. (2)].

where, in anticipation of static simulation, we have set $\alpha = 4t/r$, and where the physically irrelevant prefactor $e^{-i\alpha/4}$ and a factor of 1/4 in the exponent are included for consistency with Ref. [16].

We can go from the circuit for dynamical simulation [Eq. (4)] to the circuits needed for static quantum simulation by a VQE [2, 40, 41] by considering α as a free parameter in every occurrence of the HEIS gate. VQEs form are a promising method for computing ground state energies of various many-body systems in the NISQ era [2, 40, 41]. In VQEs, a parameterized quantum state is prepared with a parameterized circuit. The energy of this state is measured and optimized using a classical heuristic optimization routine. By the variational principle, the lowest energy that is found in this way provides an upper bound on the ground-state energy. In principle, there are no a priori restrictions on the structure of the ansatz circuit [40]. In the subclass of VQEs using the so-called Hamiltonian variational ansatz (HVA), however, each gate in the parameterized circuit is either formed by parameterized time evolution along a term in the Hamiltonian, or by parameterized time evolution according to some reference Hamiltonian [17]. The initial state of the circuit is the (known and easy-to-prepare) ground state of the reference Hamiltonian.

One possible quantum circuit following from the above considerations is depicted in Fig. 1 (colored edges). This circuit can both be used for the dynamic quantum simulation (fixed parameters α) or for static quantum simulation using the HVA (free parameters α , to be optimized by a classical optimization routine). Here, every color depicts a different layer of the circuit. In the first layer, singlet states $(|01\rangle - |10\rangle)/\sqrt{2}$ are placed along the blue edges. This provides the initial state of either the circuit for dynamical or static quantum simulation. It is the ground state of a reference Hamiltonian $\sum_{(i,j)} H_{ij}$, where the sum is over the blue edges (i, j). After preparation of the initial state, HEIS gates are placed along all orange, green, red, and blue edges. This combination of HEIS gates defines one *cucle* of the circuit. The cycle is repeated r times. (The color blue defines both the initial state and the last layer of the cycle.) Alternative circuits for dynamic or static quantum simulation (using

the HVA) may be achieved by changing the initial state or the order of the HEIS gates.

Line-graph routing maps the kagome circuit from Fig. 1 to a circuit on the heavy-hex lattice by first identifying nodes of the respective lattices, as in Fig. 1. Subsequently, Eq. (1) is applied to all gates, after which superfluous SWAP gates are removed [Eq. (2)]. The routed circuit thus obtained is shown in Fig. 2.

III. GENERAL CASE

Here, we formalize and generalize the routing strategy from the previous section, which leads to our main results, Theorem 1 and Algorithm 2. To introduce notation, we now define line- and heavy graphs more formally. The graph heavy (G) is obtained from G by defining a node i and edges (m, i), (i, m') for each edge (m, m') in G. The nodes *i* are referred to as the heavy nodes. A minor but subtle point is that this construction may lead to paths of length three as induced subgraphs of heavy(G)in which two adjacent nodes are of degree one. That is, induced subgraphs of the form $P = \{(m, i), (i, m')\},\$ with m and i nodes of degree one. For example, in the star graph (Table I), all paths emanating from the center node (i.e., the 'rays' of the star) are of the form of P. Line-graph routing [previewed in Eq. (1)] will never use mediating qubits associated with nodes locally equivalent (including neighbors and next-nearest neighbors) to node m in P and can therefore be discarded without affecting the routed circuit. In the example of the star graph (Table I), this means the qubits associated with the gray nodes may be removed from the routed circuit. To avoid frequent mention of minor but subtle point, in this paper we also refer to heavy graphs where the nodes locally equivalent (including neighbors and next-nearest neighbors) to node m in P are removed as heavy graphs. After these removals, we still refer to nodes locally equivalent (only including neighbors) to node i in P as heavy nodes.

The line graph of G, L(G), is defined in terms of heavy(G) as follows: construct heavy(G) from G and let the node set of L(G) be equal to the set of heavy nodes

 $\{i\}$ of heavy(G). An edge (i, j) is added to L(G) if and only if there exists a node m in G such that (i, m) and (m, j) are edges in heavy(G).

We say that a quantum circuit C, consisting of oneand two-qubit gates by assumption, has coupling graph G = (V, E) if V is equal to the set of qubit labels in C and if (i, j) in E if and only if there exists a twoqubit gate U_{ij} in C. If a quantum computer has coupling graph G = (V, E), then it can perform arbitrary singlequbit gates on the qubits associated with each node in Vand arbitrary two-qubit gates along each edge in E. By Eq. (1), we then have the following theorem.

Theorem 1 (Line-graph routing). Every quantum circuit C with coupling graph L(G) can be performed on quantum hardware with coupling graph heavy(G) with a SWAP overhead of at most two times the number of two-qubit gates in C.

Proof. By definition, there is a one-to-one correspondence between the nodes of L(G) and the heavy nodes $\{i\}$ of heavy(G). Therefore, the single-qubit gates of Ccan be mapped directly to hardware with coupling graph heavy(G). Furthermore, for every edge (i, j) in L(G), there are edges (i, m) and (m, j) in heavy(G), where $m = m_{ij}$ can be determined uniquely from i and j. Thus, every two-qubit gate U_{ij} in C can be mapped to the three-qubit gate $MU_{imj} := \text{SWAP}_{mi}U_{mj}\text{SWAP}_{im}$, leading to an overhead of 2λ SWAP gates, with λ the total number of two-qubit gates in C.

Note that the above theorem does not make any assumption about the graph G. Also note that the theorem provides a hierarchy of mappings, as heavy(G) itself may be a line graph. For example, invoking the above theorem twice gives a routing from L(L(G)) to heavy(L(G)).

In practice, one may be given hardware with a coupling graph H' and asked to find the class of circuits that can be run on this hardware using the line-graph construction. (For an overview of the graphs that follow, please see Fig. 3.) This task may arise if one has specific but limited quantum hardware available and wants to explore its capabilities by looking at quantum circuits that can be routed to this hardware with low overhead. The most immediate class of such quantum circuits is obviously formed by circuits with coupling graph H'. In case H' is a heavy graph, H' = heavy(G), line-graph routing extends the possibilities to the class of circuits with coupling graph L(G). If H' is a heavy graph, it is trivial to find G such that H' = heavy(G). It is also trivial to construct the line graph L(G) from G. Therefore, given a heavy hardware coupling graph, it is trivial to find the class of circuits that can be run on it by line-graph routing. For example, given hardware with a heavy-hex coupling graph, it is trivial to find that line-graph routing vields the class of kagome circuits.

The task can also be reversed: given a circuit with a coupling graph G', find the hardware on which it can be run with low overhead (cf. Fig. 3). Again, the first



Figure 3. An overview of the graphs related to line-graph routing. All maps from the right to the left are trivial given that H' is a heavy graph.

place to look would be hardware with coupling graph G'. Line-graph routing extends the possibilities by adding hardware with a coupling graph of heavy(G) with G such that G' = L(G).

But how to find G from the circuit coupling graph G'? This is less straightforward, first because it is not possible to find G from G' if G' is not a line graph. In this case, line-graph routing cannot be applied. There are numerous straightforward ways of checking whether a graph is a line graph. For example, Beineke's theorem states that a graph is a line graph if and only if it does not contain an induced subgraph out of a set of nine forbidden subgraphs [42]. One of these forbidden graphs is the claw (γ) . For example, since the heavyhex and hexagonal lattices consist entirely of claws, they are themselves not line graphs. Second, even if G' is a line graph, it is nontrivial to find the graph G such that G' = L(G). Nevertheless, Roussopoulos' algorithm [43] finds G from G', or reports that G' is not a line graph, in time $O(\max\{n, |E_{G'}|\})$, with n the number of nodes and $|E_{G'}|$ the number of edges of G'.

We briefly introduce the concepts from Roussopoulos' algorithm that are useful to us later. If G' is a line graph, Roussopoulos' algorithm partitions the edges of G' into complete subgraphs in such a way that no node lies in more than two of the subgraphs (which is possible if and only if G' is a line graph). Then, the nodes of G correspond to the sets in the partition. Additionally, the nodes that lie in only one of the sets in the partition are added as nodes of G as sets of length one. Two nodes in Gare adjacent if their corresponding sets have a nonempty overlap. For example, in Fig. 1, each triangle of colored edges forms a set in the partition of the kagome lattice because triangles are fully connected subgraphs of the kagome lattice and no node of the kagome lattice is in more than two triangles. Defining the partitions as the nodes of L(G) and putting an edge between two nodes in L(G) whenever the corresponding sets have a nonempty intersection results in the hexagonal lattice (open nodes in Fig. 1).

For line-graph routing [as previewed in Eq. (1)], L(G) must be mapped to heavy(G) in such a way that the labels of the heavy nodes of heavy(G) are identical to the labels of the nodes of L(G). The output of Roussopoulos'

algorithm provides a convenient way to achieve this labeling through the following algorithm. It takes a graph G, as generated by Roussopoulos' algorithm, as input and returns H = heavy(G), a heavy graph where the set of heavy nodes is equal to the set of nodes (i.e., with equal labels) of G' = L(G). This generalizes and automatizes the identification of nodes made in Fig. 1.

Algorithm 1 (Congruent heavy labels). Consider the edges (a, b) of G for which both a and b contain more than one node. (The nodes in a and b are from G', the latter of which need not be provided explicitly to the current algorithm.) For all such edges (a, b), add (a, c) and (c, b) to an empty graph H. Because a and c are distinct sets in a partition of the edges of G' into complete subgraphs and a and c have a nonempty overlap, c contains a single entry and this entry is a node from G'. Now, consider the edge cases, that is, the edges (a, b) in G where either a or b is a set of length one. For every such edge, add (a, b) to H.

Everything is now in place for the succinct presentation of line-graph routing. It maps any circuit C, with unknown coupling graph G', to a circuit with coupling graph H whenever G' is a line graph, G' = L(G), and where H = heavy(G).

Algorithm 2 (Line-graph routing). Construct the coupling graph G' of the circuit C. Run Roussopoulos' algorithm on G' to obtain $G = L^{-1}(G')$. Construct heavy(G) using Algorithm 1. For every two-qubit gate U_{ij} in C, let $m = m_{ij}$ be the node in heavy(G) that is in between nodes i, j of heavy(G) and replace U_{ij} by $MU_{imj} := \text{SWAP}_{mj}U_{im}\text{SWAP}_{jm}$.

Note the existence and uniqueness of mediating qubit m are guaranteed by the definition of line graphs. In some cases, some qubits, related to the so-called lone leaf nodes in the coupling graph of the output circuit of line-graph routing, can be removed from that circuit, reducing the SWAP and qubit count of line-graph routed circuits. This removal leads to a marginal improvement, which may nevertheless be crucial given hardware with few qubits, but arguably obfuscates the general idea of line-graph routing. The definition of lone leaf nodes and the method for their removal is presented in Appendix A.

As opposed to heuristic methods [7], line-graph routing is deterministic, allowing rigorous performance guarantees. Given the time complexity of Roussopoulos' algorithm [43], it is straightforward to show that the time complexity of line-graph routing is $O(\Lambda^2)$, with Λ the number of gates in the input circuit C. A tighter bound on the time complexity can possibly obtained, but this requires a more sophisticated analysis. Such an analysis is unnecessary for the current purposes because a nonoptimized implementation of line-graph routing already routes circuits with thousands of qubits and hundred thousands of gates within a minute on a standard laptop [44].

To utilize line-graph routing in practice, one additional step may be required. Quantum hardware providers will generally use a specific labeling of the qubits on their hardware, leading to a hardware coupling graph H'. This labeling may differ from the labeling of H = heavy(G)obtained through Algorithm 1. However, H can be mapped to a subgraph of H' using an algorithm that finds subgraph isomorphisms, such as the VF2 algorithm [45]. The VF2 algorithm generates a list of isomorphic subgraphs, and in practice, a selection is made based on a performance metric, such as the average two-qubit gate fidelity, to identify the subgraph with the best performance. This step is commonly implemented in quantum computing software development kits. In this paper, we also refer to H as the hardware coupling graph and in this wording leave the possible mapping to the fixed qubit labels provided by a hardware provider implicit.

IV. IMPLEMENTATION AND BENCHMARKING

In the Supplemental Material [44], we implement, showcase and benchmark line-graph routing, Algorithm 2, together with the removal of lone leafs (Appendix A). The implementation takes any Qiskit [46] quantum circuit consisting of one- and two-qubit gates, constructs its coupling graph L(G) or reports that the coupling graph is not a line graph, finds G and heavy(G), and outputs the routed circuit with coupling graph heavy(G). The implementation does not rely heavily on Qiskit's methods and can hence straightforwardly be transformed to an implementation in other quantum software development kits.

Line-graph routing is benchmarked against all routing methods available in Qiskit by default [44, 46]. In this section, we show line-graph routing is able to confidently outperform these default methods on relevant problem instances. There are also problem instances where linegraph routing does not outperform the default methods. In the end of this section, we discuss for which types of instances we expect line-graph routing to be superior. We consider two types of circuits.

(i) Random. With probability 2/5, a CNOT gate is placed along a randomly chosen edge of a given virtual graph. With a probability 3/5 a gate from the set $\{H, S, T\}$ is chosen uniformly at random and placed at a random node.

(ii) Quantum simulation. As described in detail for the kagome lattice (Sec. II B), circuits for the dynamical and static quantum simulation of the HAFM on any lattice can be defined by an edge-coloring of that lattice [15, 16]. An edge coloring of a graph is an assignment of colors to the edges such that no two edges with the same color are incident on the same node. This edge coloring is called minimal if it uses the least possible number of colors. We perform edge coloring of the virtual lattices by an automatic method that generally finds an edge coloring

L(G)	Routing method	Opt. depth	$n_{\rm SWAP}$	$n_{ m qubit}$	$t_{\rm tot}~({\rm s})$	\bar{t} (s)
Complete	line-graph	1036	650	10	<1	<1
	SABRE	720	146	9	3	<1
Kagome	line-graph	226	7968	300	43	43
	SABRE	790	8486	200	886	55 ± 1
Shuriken	line-graph	209	13600	476	77	77
	SABRE	1099	17064	323	2520	157 ± 4
Checkerboard	line-graph	435	18521	393	67	67
	SABRE	2121	23060	282	1750	109 ± 1

Table II. Excerpt of the benchmarking data available in the Supplemental Material [44]. The column headers are defined in the main text.

that is not minimal. The benefit of this coloring method is that it does not require a manual assignment of edge colors. The downside is that we expect line-graph routing to work best (compared to other methods) for circuits derived from a minimal edge coloring. This does not pose a problem because, as we show in this section, line-graph routing is already able to outperform the default methods in Qiskit in routing circuits derived from a nonminimal edge coloring.

We found SABRE [7] to outperform the other methods in Qiskit and therefore we focus on a comparison between line-graph routing and SABRE in what follows. Unlike line-graph routing, SABRE is a probabilistic routing method that obtains a different qubit routing with each run. Additionally, the intensity of the optimization that is part of SABRE can be varied, leading to a tradeoff between the classical resources required and the performance characteristics of the routed circuits. We address these issues by running SABRE 16 times (at fixed optimization level) and comparing the performance against one run of line-graph routing. We do this separately for every optimization level available by default in Qiskit, which range from optimization level 0 ('no optimization') to optimization level 3 ('heavy-weight optimization') in integer steps [46].

In Table II, we show an excerpt of the benchmarking data, which includes problem instances on which linegraph routing does and does not perform well. The following performance characteristics are listed.

(i) Opt. depth. The optimal (lowest) depth reached by the routing method among all runs (line-graph routing is run once per virtual graph, SABRE is run 16 times per virtual graph). Routed circuits are obtained by inserting SWAP gates (as dictated by line-graph routing or SABRE) and no further gate identities are used to simplify the resulting circuits. So, in the case of random input circuits, the routed circuits consist of gates from the set {CNOT, H, S, T} \cup {SWAP}. The routed circuits contain, for example, double H gates if those were present in the input circuit. In the case of quantum simulation input circuits, the routed circuits consist of gates from the set {SINGLET, HEIS(α)} \cup {SWAP}.

(ii) n_{SWAP} . The number of SWAP gates of the routed

circuit that achieved the lowest depth.

(*iii*) n_{qubit} . The number of active qubits in the routed circuit that achieved the lowest depth.

 $(iv) t_{tot}$. The total wall-clock time needed to run all runs of the routing method. For line-graph routing, this includes the time needed to find heavy(G) from L(G). The implementation of line-graph routing repeatedly loops through all gates using (slow) Python loops and can likely be sped up considerably, if needed. We use Qiskit's standard implementation of SABRE. SABRE is given the target graph heavy(G) as input and hence finding heavy(G) from L(G) is not included in its wall-clock time. The benchmarks for different methods are always run on the same hardware.

 $(v) \ \bar{t}$. The average wall-clock time for a single run of the routing method. The error bars are calculated by bootstrapping the individual wall-clock times and represent symmetrized 95% confidence intervals.

The first two data lines of Table II show the performance characteristics of line-graph routing and SABRE when applied to a random circuit on the complete graph with 9 nodes. SABRE was run with optimization level 1. (Passing a higher optimization level to Qiskit's transpiler will trigger the usage of gate identities.) Already at an optimization level of 1 SABRE outperforms line-graph routing on all performance characteristics considered except the total wall-clock time.

Data lines 3 and 4 of Table II show the performance characteristics of line-graph routing and SABRE when applied to circuits for the quantum simulation of the HAFM on patches of the kagome lattice measuring 7×7 unit cells, with open boundary conditions and padded edges (see [44]). Here, SABRE was run at Qiskit's transpiler optimization level 3, the highest optimization level available. As opposed to the transpilation of the random circuits, no gate identities are used in this process because the gates {SWAP, SINGLET, HEIS} have unknown properties to Qiskit's transpiler [46]. Line-graph routing outperforms SABRE in terms of the optimal depth of the routed circuits by about a factor of 3.5, while at the same time requiring less time than one run of SABRE. The line-graph routed circuit also uses less SWAP gates than the SABRE routed circuit with the lowest depth.

However, the line-graph routed circuit uses a factor of 1.5 more qubits than the SABRE routed circuit that achieved the lowest depth. Nevertheless, the space-time volume of the line-graph routed circuit is still about half of the space-time volume of the SABRE routed circuit with the lowest depth. Similar results, with an even larger performance gap, hold for the routing of circuits for the quantum simulation of the HAFM on patches of the shuriken lattice measuring 7×7 unit cells (lines 5 and 6 of Table II) and patches of the checkerboard lattice of 7.5×7.5 unit cells (last two lines of Table II).

Line-graph routing was conceived while keeping in mind its application in mapping quantum circuits to hardware with a lattice-like low-degree coupling graph. It is therefore not expected to perform well in mapping quantum circuits to hardware with a high-degree coupling graph without lattice-like structure. In fact, the benchmarking results show that line-graph routing is not well-suited for mapping circuits on the complete coupling graph to hardware with a star coupling graph. One property of line-graph routing that leads to its low effectivity on unstructured, high-degree graphs is that in the output circuit of line-graph routing, the qubits are assigned a base location where they return to after they are acted on by one [Eq. (1)] or multiple [Eq. (2)] gates from the input circuit. This is likely advantageous for input quantum circuits with a structured, low-degree coupling graph. For example, in Fig. 2, in the third layer of the cycle (red gates), all qubits are still close to where they are needed, leading to the insertion of few SWAP gates to get them there.

However, on high-degree coupling graphs, the property of a base location need not be advantageous since any qubit can be routed to any other qubit in relatively few steps and regularly returning qubits to their base location leads to the insertion of unnecessary SWAP gates. As an extreme example, let us look at the action of linegraph routing on a circuit C with a complete coupling graph, where at one point in C a cascade of CNOT gates, $\prod_{i=n}^{1} \text{CNOT}_{i,i+1}$, is prescribed. Line-graph routing maps the circuit C to a circuit on the star graph (Table I). To perform $CNOT_{i,i+1}$ on star-graph hardware, line-graph routing first swaps qubit i to the center of the star, performs a CNOT between the center qubit and qubit i+1, and swaps qubit i back to its original position [Eq. (1)]. Not insisting that the qubits eventually return to their original position leads to the possibility of more efficient routing strategy. To perform the $CNOT_{i,i+1}$, swap qubit i to the center, perform the CNOT between the center qubit and qubit i + 1, and leave qubit i in the center. After repeating this procedure for the subsequent CNOT in the circuit, qubit *i* will end up at the initial location of qubit i+1. The latter approach uses half of the number of SWAP gates compared to line-graph routing.

A second situation in which line-graph routing is not expected to perform well is when there are subgraphs of the virtual graph L(G) that are isomorphic to subgraphs of heavy(G). In that case, unnecessary mediating qubits may be inserted. For example, let L(G) be the path graph on n nodes. When applied to circuits on this graph, line-graph routing introduces mediating qubits between all qubits of L(G), leading to a circuit on a path graph with n-1 added mediating qubits. This happens despite the fact that no routing is needed at all. If needed, linegraph routing can be enhanced such that it detects and prevents this behavior.

V. DISCUSSION AND OUTLOOK

In this paper, we developed line-graph routing, a qubit routing strategy that efficiently and deterministically maps any quantum circuit on a line graph L(G) to a circuit on the heavy graph heavy(G). By software implementation and benchmarking, we showed its ability to outperform standard, general-purpose methods on input quantum circuits, circuit sizes, and hardware connectivity graphs of practical relevance.

Line-graph routing showed not to perform well in mapping circuits on the complete graph to the star graph. We attribute this to the high degree and absence of a lattice-like structure of the complete graph. Based on our benchmarking results, we expect line-graph routing to outperform general-purpose methods in routing circuits with a line-graph coupling graph to low-degree hardware coupling graphs. For superconducting qubits, these are exactly the hardware coupling graphs preferable from an engineering standpoint [31, 32, 36].

Line-graph routing is limited in its applicability because it is only defined for pairs of graphs (L(G), heavy(G)). General-purpose methods are able to map any circuit on any graph to a circuit on any other graph (given that the latter graph is connected and has at least the same number of nodes as the former graph). However, in general, finding the optimal routing strategy is NP-hard. It is therefore unlikely that general-purpose methods can find the optimal or closeto-optimal routing strategy. Routing strategies that are defined on a subset of possible input circuits can nevertheless be highly efficient and effective, and indeed linegraph routing is shown to have the ability to outperform standard-purpose methods on the subset of problem instances for which it is defined. The set of pairs of graphs (L(G), heavy(G)) is still infinitely large and contains pairs of practical relevance.

The routed circuit for the simulation of a nearestneighbor model on L(G) can be reinterpreted as a circuit for the quantum simulation of a next-nearest-neighbor model on heavy(G). This is evident from Table I and Fig. 1 and follows directly from the definition of line graphs as given in Sec. III. This routed circuit for the simulation of a nearest-neighbor model on L(G) can be lifted to a circuit for the quantum simulation of a model containing both nearest- and next-nearest-neighbor interactions on heavy(G) by adding the gates arising from nearest-neighbor interactions on heavy(G) to the routed circuit. These gates naturally satisfy the heavy(G) connectivity graph and can thus be added to the circuit without any routing. This further enlarges the range of applicability of line-graph routing.

The effectiveness of line-graph routing may be improved further by leveraging the freedom of which qubit is swapped with the mediating qubit. In line-graph routing, for the implementation of the gate U_{ij}^{ℓ} , it is qubit *i* that is swapped with the mediating qubit [Eq. (1)]. However, in some cases, it is only when swapping qubit *j* with the mediating qubit that a cancellation of SWAP gates from consecutive mediated two-qubit gates MU [Eq. (1), Fig. 2] occurs. Thus, line-graph routing may be improved by letting the decision of which qubit to swap with the mediating qubit depend on the ensuing gates in the input quantum circuit.

Although we showed the ability of line-graph routing to outperform general-purpose methods, we did not prove it gives the optimal routing. In fact, it was shown to be suboptimal in cases where the input circuit is a circuit on the complete graph. A proof or refutation for the optimality of line-graph routing (possibly after including the improvement of the previous paragraph) would therefore require careful consideration of the allowed input circuits and performance characteristic for which optimality is considered. Such a proof may inspire even A straightforward but exciting way to carry on the work in the current paper is to run our routed circuits on actual quantum hardware. The circuits available in the Supplemental Material [44] can be executed as-is on hardware with the appropriate hardware coupling graphs. Remaining challenges therein are the extraction of useful physical quantities from the generated quantum states. To obtain results that go beyond anything that can be obtained classically requires further improvements of error mitigation techniques and quantum hardware.

All code and data used to generate the results in this paper are available as Supplemental Material [44] and at Ref. [47].

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Appendix A: Removal of lone leaf nodes

In some cases, line-graph routing produces circuits containing qubits i that are swapped with only one corresponding mediating qubit m_i during the course of the entire circuit. In these cases, the mediating qubits m_i can be removed and replaced by the corresponding qubits i, leading to a reduction in the qubit and SWAP count of the routed circuit.

We define a *lone leaf* i as a node in a coupling graph heavy(G) that is of degree one and shares its only edge with a node that has no other neighbors of degree one. We do not assume the coupling graph is a tree. As an example, consider the triangle graph with an extra node i connected to one of the nodes m_i of the triangle. The node i is a lone leaf.

If a node i is a lone leaf in a coupling graph heavy(G),

by the construction of line-graph routing its neighbor mmust be a mediating node. For every two-qubit gate that is performed between qubit i and any other neighbor jof m, qubits i and m first need to be swapped. So, as far as the interactions with qubit i are concerned, qubit m may be fully eliminated by simply removing qubit mand reverting mediated two-qubit gates MU_{imj} to the original gates MU_{imj} . Physically, qubit i can be put on the place of qubit m. After the removal of m and the relocation of qubit i, inspection of Eq. (1) reveals that qubit i may in fact act as a mediating node for any twoqubit gates between any two neighbors j, k of m unequal to i; after a mediated gate MU_{jik} , qubit i returns to its starting position unaffected.

Supplemental material to: "Line-graph qubit routing: from kagome to heavy-hex and more" Joris Kattemölle¹ and Seenivasan Hariharan²

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Abstract

We showcase and benchmark line-graph qubit routing by routing random circuits (with fixed virtual graph) and circuits needed for the quantum simulation of the Heisenberg antiferromagnet on various graphs. The code implementing line-routing itself is found in line_graph_routing.py and maps Qiskit quantum circuits to Qiskit quantum circuits. (This can be altered to other circuit libraries with little effort.) We benchmark our results against other methods. The interactive version of this document is line_graph_routing.ipynb.

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1 Requirements

This notebook should typically run after installing the following packages with pip (or conda). In a terminal, run

pip install qiskit[visualization]

or

pip install 'qiskit[visualization]'

and

pip install netket networkx tabulate

Note Netket currently needs Python 3.9 (and SciPy >= 1.9.3). Netket is only used to generate patches of the kagome lattice as graphs and not for linegraph routing itself. This notebook was tested with a pip environment that can be recreated with requirements.txt by running pip install -r requirements.txt (after creating a new environment).

The file line_graph_routing.py should be placed in the same folder as the current notebook.

```
import line_graph_routing as lgr
import networkx as nx
```

2 Kagome to heavy-hex

2.1 Random

Create a random circuit on a patch of the kagome lattice of 3×3 unit cells and show the circuit's coupling graph. With probability 2/5, a CNOT is placed along an edge of the connectivity graph. With a probability 3/5 a gate from {H,S,T} is chosen uniformly at random and placed at a random node.

```
lg = lgr.kagome(3, 3)
qc = lgr.random_circuit(lg, 10**4)
cg = lgr.coupling_graph(qc)
nx.draw_kamada_kawai(cg)
print(qc.depth())
```



Route the circuit to a circuit with heavy-hex coupling graph.

```
qc = lgr.line_graph_route(qc)
cg = lgr.coupling_graph(qc)
nx.draw_kamada_kawai(cg)
print(qc.depth())
```

1337



2.2 Quantum simulation

We base our circuits on edge colorings of the (kagome) lattice by identifying every color with a layer of HEIS gates. One of the colors (color '0') doubles as a color specifying the initial state by indicating along which edges singlet states are placed. The entire circuit is repeated p times, excluding initial state preparation. Every HEIS gets its own parameter. These parameters can later be bound to specific values to obtain circuits for dynamical quantum simulation or for simulated adiabatic state preparation.

First, create and show an edge coloring of the kagome lattice.

```
lg = lgr.edge_coloring(lg)
lgr.draw_edge_coloring(lg)
```

Matching is perfect Edge coloring is not minimal



Create the associated circuit, route it to heavy-hex hardware and show the coupling graph of the routed circuit.

```
p = 1
qc = lgr.heis_circuit(lg, p)
print(qc.depth())
qc = lgr.line_graph_route(qc)
print(qc.depth())
cg = lgr.coupling_graph(qc)
nx.draw_kamada_kawai(cg)
```

6

```
15
```



Since the above circuit has the connectivity of a 2D graph, the corresponding quantum circuit diagram will not be very insightful. However, for a single unit cell of the kagome lattice, the routed quantum circuit becomes a circuit on a circle, which allows for a clear representation as a quantum circuit diagram. We first

create an edge coloring of the unit cell patch.

lg = lgr.kagome(1, 1)
lg = lgr.edge_coloring(lg)
lgr.draw_edge_coloring(lg)





Create the HEIS circuit ased on this coloring, map it to heavy-hex hardware, and show the coupling graph of the routed circuit.

4 7



Show the circuit diagram of the routed circuit, with parameters al_i.



The fircuit depth can be reduced further by replacement of the initial and final SWAP gates between qubits (10,7) and (8,6) by a relabeling of those qubits.

3 Complete graph to star graph

3.1 Random

Create a random circuit on the complete graph of four nodes and show the circuit's coupling graph.

```
n = 5
lg = nx.complete_graph(4)
qc = lgr.random_circuit(lg, 10**2)
cg = lgr.coupling_graph(qc)
nx.draw_kamada_kawai(cg)
print(qc.depth())
```

56



Route the circuit to a circuit with star-graph connectivity.

```
qc = lgr.line_graph_route(qc)
cg = lgr.coupling_graph(qc)
nx.draw_kamada_kawai(cg)
```

print(qc.depth())





3.2 Quantum simulation

As before, circuits are defined by identifying every color with a layer of HEIS-gates. For more details, see the kagome to heavy-hex section.

Create and show an edge coloring of the complete graph.

```
lg = lgr.edge_coloring(lg)
lgr.draw_edge_coloring(lg)
```

Matching is perfect Edge coloring is not minimal



Create the associated circuit, route it to heavy-hex hardware, and show the coupling graph of the routed circuit.

```
p = 1
qc = lgr.heis_circuit(lg, p)
print(qc.depth())
qc = lgr.line_graph_route(qc)
```



We do not show the circuit diagram in this case because the routed circuit is not a circuit on a line.

4 Shuriken to heavy squareoctagon

4.1 Random

Create a random circuit on a patch of the shuriken lattice of 3×3 unit cells.

```
n = 5
lg = lgr.shuriken(3, 3)
qc = lgr.random_circuit(lg, 10**4)
cg = lgr.coupling_graph(qc)
nx.draw_kamada_kawai(cg)
print(qc.depth())
```

517



Route the circuit to a circuit with heavy-squareoctagon connectivity.

```
qc = lgr.line_graph_route(qc)
cg = lgr.coupling_graph(qc)
nx.draw_kamada_kawai(cg)
print(qc.depth())
```

877



4.2 Quantum simulation

As before, circuits are defined by identifying every color with a layer of HEIS-gates. For more details, see the kagome to heavy-hex section.

Create and show an edge coloring of the shuriken lattice.

```
lg = lgr.edge_coloring(lg)
lgr.draw_edge_coloring(lg)
```

Matching is perfect Edge coloring is not minimal



Create the associated circuit, route it to heavysquare-octagon hardware, and show the coupling graph of the routed circuit.

```
p = 1
qc = lgr.heis_circuit(lg, p)
print(qc.depth())
qc = lgr.line_graph_route(qc)
print(qc.depth())
cg = lgr.coupling_graph(qc)
nx.draw_kamada_kawai(cg)
```

6 12



Again, the resulting circuit diagram will not be very insightful, but it will be for a single-unit cell patch of the shuriken lattice.

```
lg = lgr.shuriken(1, 1)
lg = lgr.edge_coloring(lg)
lgr.draw_edge_coloring(lg)
```

Matching is perfect Edge coloring is minimal



p = 1 qc = lgr.heis_circuit(lg, p) print(qc.depth()) qc = lgr.line_graph_route(qc)



qc.draw('latex')

5 9



5 Checkerboard to heavysquare

m = 2.5 # For the checkerboard lattice,
→ specify dimentions in nodes by nodes.
lg = lgr.checkerboard(m, m)
qc = lgr.random_circuit(lg, 10**4)
cg = lgr.coupling_graph(qc)
nx.draw_kamada_kawai(cg)
print(qc.depth())

872



Route the circuit to a circuit with a heavy-square coupling graph.

```
qc = lgr.line_graph_route(qc)
cg = lgr.coupling_graph(qc)
nx.draw_kamada_kawai(cg)
print(qc.depth())
```

1622



5.1 Quantum simulation

As before, circuits are defined by identifying every color with a layer of HEIS-gates. For more details, see the kagome to heavy-hex section.

Create and show an edge coloring of the checkerboard ~~6 lattice

```
lg = lgr.edge_coloring(lg)
lgr.draw_edge_coloring(lg, spectral=True)
→ # Use spactral method to find location
→ of nodes.
```

Matching is perfect Edge coloring is not minimal



Create the associated circuit, route it to heavysquare-octagon hardware, and show the coupling graph of the routed circuit.

p = 1 qc = lgr.heis_circuit(lg, p) print(qc.depth()) qc = lgr.line_graph_route(qc) print(qc.depth()) cg = lgr.coupling_graph(qc) nx.draw_kamada_kawai(cg)

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6 Random line graph to random heavy graph

6.1 Random

Create a random circuit on a random graph with 6 nodes and show the circuit's coupling graph. For details on radom_line_graph generation, see its function definition in line_graph_routing.py.

```
n = 6
lg = lgr.random_line_graph(6)
qc = lgr.random_circuit(lg, 10**3)
cg = lgr.coupling_graph(qc)
nx.draw_kamada_kawai(cg)
print(qc.depth())
```

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Route the circuit to a circuit with the associated heavy connectivity.

```
qc = lgr.line_graph_route(qc)
cg = lgr.coupling_graph(qc)
nx.draw_kamada_kawai(cg)
print(qc.depth())
```

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6.2 Quantum simulation

As before, circuits are defined by identifying every color with a layer of HEIS-gates. For more details, see the kagome to heavy-hex section.

Create and show an edge coloring of random graph. The method we use to find a perfect matching (needed for initial state preparation) is limited and may not find a perfect matching even if it exists. If a perfect matching is not found, try to create another random line graph (i.e., evaluate the two cells above) or use more sophisticated (or brute-force) methods to find a perfect matching.

lg = lgr.edge_coloring(lg)
lgr.draw_edge_coloring(lg)

Matching is perfect Edge coloring is not minimal



Create the associated circuit, route it to heavy-hex hardware and show the coupling graph of the routed circuit.

p = 1 qc = lgr.heis_circuit(lg, p) print(qc.depth()) qc = lgr.line_graph_route(qc) print(qc.depth()) cg = lgr.coupling_graph(qc) nx.draw_kamada_kawai(cg)





We do not show the circuit diagram in this case because the routed circuit is not a circuit on a line.

7 Benchmarking

We benchmark line graph routing by performing the above routing tasks (but for larger unit cells) using both line-graph qubit routing and all methods available in qiskit. These methods are 'basic', 'lookahead', 'stochastic', and 'sabre' [1].

The benchmarking settings are specified by the following options: - name The name of the virtual graph, either kagome, shuriken or complete. - size. In case of kagome and shuriken: the size of the patch in unit cells by unit cells. In case of complete: the number of nodes of the complete graph. - circuit type. Either quantum simulation or random, as presented in this notebook. - p In case of kagome and shuriken: the number of cycles of the circuit. In case of complete: the number of random gates from the set H,T,S,CNOT. - repetitions. The number of runs for the methods sabre and stochastic. The methods line-graph and basic are deterministic and hence only run once. Correspondingly, the reported total time pertains to the time taken for this single run in case of the latter two methods. - optimization_level. Either 0, 1, 2, or 3. This specifies the optimization level used for the routing methods implemented in qiskit [2]. This parameter is passed directly to Qiskit's transpiler [1]. - methods. The methods to benchmark line graph transpilation against. Must be a list containing elements from ['sabre','basic','lookahead',stochastic]. These methods are passed directly to Qiskit's transpiler [1].

The methods **sabre** and **stochastic** are probabilistic, achieving a different routing each time they are run, and hence achieve different performance characteristics with each run. We therefore run these methods *repetitions* times and report the average, confidence interval, and best performance out of these runs. Error bars on the data show the (symmetrized) 95% confidence interval and are obtained by bootstrapping the data. The error interval for num_qubits is sometimes given by nan because in those cases the number of qubits was equal for all runs. The routing methodsline-graph and basic are deterministic and for these we enforce repetitions=1.

We consider the following performance characteristics. - method The routing method. - av. n_swaps The average number of swaps obtained among the repetitions runs of the routing method. - min n_swap The number of swaps of the run that achieved the lowest depth. - av. depth The average depth of the routed circuits among the repetitions runs of the routing method. We focus on the performance of routing so none of the gates in any of the routing methods are compiled into hardware native gates. That is, for the purposes of assessing routing performance, we assume the gate set SWAP, HEIS, H, X, Z CNOT for the quantum simulation circuits. For the random circuits we assume SWAP, CNOT, H, S, T. - min depth The minimum depth among the repetitions runs of the routing method. - av. n_qubits The average number obtained among the repetitions runs of the routing method. - av. time The average (minimum) wall clock time of the routing runs. - min. time The number wall clock run time of the run that achieved the lowest depth.

[1] Qiskit 0.43.0 documentation, https://qiskit.org/documentation/stubs/qiskit.compiler.transpile.html, accessed 11h May 2023.

[2] https://github.com/Qiskit/qiskit-terra/tree/main/qiskit/transpiler/preset_passmanagers

7.0.1 Quantum simulation, kagome and shuriken, agianst SABRE

```
import line_graph_routing as lgr # Loading these makes these cells stand-alone
import pickle
settings=[]
for name in ['kagome', 'shuriken']:
    for side in range(1,9,2):
        for p in [1,8,16]:
            for optimization_level in range(4):
               setting={'name':name,
                    'size': (side,side),
```

```
'circuit_type': 'quantum_simulation',
                     'p': p,
                     'repetitions' : 16,
                     'optimization_level' : optimization_level,
                     'methods' : ['sabre']
                    }
                settings.append(setting)
## Uncomment to rerun benchmarks. This takes a couple of hours.
#results=[]
#for setting in settings:
   result=lgr.benchmark(**setting)
#
#
  results.append(result)
    lgr.print_benchmark(result)
#
#
#with open('benchmark_data/kagome_shuriken.pkl', 'wb') as f:
  pickle.dump(results,f)
#
#Load previously obtained results from disk and show them.
import pickle
with open('benchmark_data/kagome_shuriken.pkl','rb') as f:
   results=pickle.load(f)
for result in results:
```

```
lgr.print_benchmark(result)
```

name = kagome, size = (1, 1), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 0											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	12 ± 0.0	12	7 ± 0.0	7	12 ± 0.0	12	0.08	0.08 ± 0.0	0.08		
sabre	16 ± 2.06	6	12.38 ± 2.09	6	10.25 ± 1.0	12	0.7	0.04 ± 0.06	0.02		

name = kagome, size = (1, 1), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 1										
mothod		min n arran	ou donth	min donth	ov navbita	min auhita	total time (a)	av time (a)	min time (a)	
mecnou	av. n_swaps	min. n_swap	av. depth	min. depen	av. n_qubits	min. qubits	COLAI CIME (S)	av. time (s)	min. time (s)	
line graph	12 + 0.0	10	7 + 0 0	7	12 + 0.0	10	0.08	0.08 + 0.0	0.08	
11ne-graph	12 ± 0.0	12	7 ± 0.0	'	12 ± 0.0	12	0.08	0.08 ± 0.0	0.08	
sabre	10 ± 0.81	6	6.44 ± 0.34	6	8.75 ± 0.75	8	0.32	0.02 ± nan	0.02	

<pre>name = kagome, size = (1, 1), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 2</pre>											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	12 ± 0.0	12	7 ± 0.0	7	12 ± 0.0	12	0.08	0.08 ± 0.0	0.08		
sabre	6 ± nan	6	6.0 ± nan	6	8.0 ± nan	8	0.33	0.02 ± 0.0	0.02		

name = kago	ame = kagome, size = (1, 1), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 3										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	12 ± 0.0	12	7 ± 0.0	7	12 ± 0.0	12	0.08	0.08 ± 0.0	0.08		
sabre	6 ± nan	6	6.0 ± nan	6	8.0 ± nan	8	0.75	0.05 ± 0.02	0.04		

name = kagon	me, size = (1, 1), circuit_typ	e = quantum_si	mulation, p =	8, repetitions =	16, optimizatio	on_level = 0		
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	96 ± 0.0	96	49 ± 0.0	49	12 ± 0.0	12	0.67	0.67 ± 0.0	0.67
sabre	99 ± 7.66	48	52.94 ± 5.28	41	11.5 ± 0.75	12	1.39	0.09 ± 0.02	0.08

name = kago	me, size = (1,	 circuit_typ 	e = quantum_s:	imulation, p =	8, repetitions :	= 16, optimizati	on_level = 1		
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	96 ± 0.0	96	49 ± 0.0	49	12 ± 0.0	12	0.64	0.64 ± 0.0	0.64
0F									
sabre	48 ± 8.62	48	42.88 ± 3.28	41	8.5 ± 0.75	8	1.67	0.1 ± 0.02	0.09

name = kagome, size = (1, 1), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 2										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	96 ± 0.0	96	49 ± 0.0	49	12 ± 0.0	12	0.64	0.64 ± 0.0	0.64	
sabre	48 ± 7.19	48	41.56 ± 1.41	41	8.25 ± 0.62	8	2.07	0.13 ± 0.02	0.11	

<pre>name = kagome, size = (1, 1), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 3</pre>											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	96 ± 0.0	96	49 ± 0.0	49	12 ± 0.0	12	0.64	0.64 ± 0.0	0.64		
sabre	48 ± nan	48	41.0 ± nan	41	8.0 ± nan	8	5.07	0.32 ± 0.02	0.29		

name = kagome, size = (1, 1), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 0											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	192 ± 0.0	192	97 ± 0.0	97	12 ± 0.0	12	1.29	1.29 ± 0.0	1.29		
sabre	100 ± 18.19	96	99.38 ± 9.69	81	11.25 ± 0.75	12	2.78	0.17 ± 0.02	0.15		

name = kago	name = kagome, size = (1, 1), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 1										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	192 ± 0.0	192	97 ± 0.0	97	12 ± 0.0	12	1.31	1.31 ± 0.0	1.31		
sabre	96 ± 17.62	96	86.69 ± 8.53	81	8.5 ± 0.75	8	3.18	0.2 ± 0.03	0.17		
name = kago	me, size = (1,	 circuit_typ 	e = quantum_si	mulation, p =	16, repetitions	= 16, optimizat	ion_level = 2				

method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	192 ± 0.0	192	97 ± 0.0	97	12 ± 0.0	12	1.3	1.3 ± 0.0	1.3
sabre	96 ± nan	96	81.0 ± nan	81	8.0 ± nan	8	3.74	0.23 ± 0.02	0.21

name = kagome, size = (1, 1), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 3										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	192 ± 0.0	192	97 ± 0.0	97	12 ± 0.0	12	1.29	1.29 ± 0.0	1.29	
sabre	96 ± nan	96	81.0 ± nan	81	8.0 ± nan	8	10.26	0.64 ± 0.02	0.67	

<pre>name = kagome, size = (3, 3), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 0</pre>										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	112 ± 0.0	112	15 ± 0.0	15	68 ± 0.0	68	0.6	0.6 ± 0.0	0.6	
sabre	307 ± 3.44	286	72.62 ± 3.5	63	55.0 ± 1.66	55	1.83	0.11 ± 0.02	0.09	

name = kagome, size = (3, 3), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 1											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	112 ± 0.0	112	15 ± 0.0	15	68 ± 0.0	68	0.6	0.6 ± 0.0	0.6		
sabre	85 ± 3.56	85	27.38 ± 1.34	22	45.25 ± 1.66	52	1.8	0.11 ± 0.02	0.1		

name = kagome, size = (3, 3), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 2										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	112 ± 0.0	112	15 ± 0.0	15	68 ± 0.0	68	0.6	0.6 ± 0.0	0.6	
sabre	96 ± 4.19	82	26.81 ± 2.41	20	46.19 ± 1.12	47	2.62	0.16 ± 0.02	0.15	

name = kagome, size = (3, 3), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 3										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	112 ± 0.0	112	15 ± 0.0	15	68 ± 0.0	68	0.5	0.5 ± 0.0	0.5	
sabre	64 ± 2.34	67	21.44 ± 1.94	14	46.0 ± 1.75	47	7.77	0.49 ± 0.02	0.55	

name = kagome, size = (3, 3), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 0										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	896 ± 0.0	896	113 ± 0.0	113	68 ± 0.0	68	4.25	4.25 ± 0.0	4.25	
sabre	1361 ± 33.51	1203	312.5 ± 9.31	287	59.5 ± 1.75	51	10.83	0.68 ± 0.03	0.69	

name = kagome, size = (3, 3), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 1										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	896 ± 0.0	896	113 ± 0.0	113	68 ± 0.0	68	4.39	4.39 ± 0.0	4.39	
sabre	625 ± 19.67	613	192.38 ± 14.0	149	45.19 ± 1.28	44	10.99	0.69 ± 0.04	0.76	

name = kagome, size = (3, 3), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 2										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	896 ± 0.0	896	113 ± 0.0	113	68 ± 0.0	68	4.42	4.42 ± 0.0	4.42	
sabre	621 ± 9.75	621	182.06 ± 9.79	152	43.56 ± 1.4	44	14.53	0.91 ± 0.04	0.95	

name = kagome, size = (3, 3), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 3										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	896 ± 0.0	896	113 ± 0.0	113	68 ± 0.0	68	4.23	4.23 ± 0.0	4.23	
sabre	601 ± 10.06	555	176.25 ± 7.62	149	43.44 ± 0.72	43	45.11	2.82 ± 0.04	2.74	

name = kagome, size = (3, 3), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 0										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	1792 ± 0.0	1792	225 ± 0.0	225	68 ± 0.0	68	8.45	8.45 ± 0.0	8.45	
sabre	2627 ± 55.91	2304	553.94 ± 12.78	509	60.94 ± 2.12	63	20.65	1.29 ± 0.05	1.29	

name = kagome, size = (3	, 3), circuit_type = quantum_simulation, p	= 16, repetitions = 16, optimization_	level = 1	

method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	1792 ± 0.0	1792	225 ± 0.0	225	68 ± 0.0	68	8.87	8.87 ± 0.0	8.87
sabre	1213 ± 28.0	1213	388.12 ± 18.66	318	43.75 ± 1.16	44	20.66	1.29 ± 0.05	1.35

name = kagome, size = (3, 3), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 2									
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	1792 ± 0.0	1792	225 ± 0.0	225	68 ± 0.0	68	8.4	8.4 ± 0.0	8.4
sabre	1215 ± 14.41	1142	377.69 ± 19.52	311	44.06 ± 1.19	49	28.01	1.75 ± 0.01	1.79

name = kagome, size = (3, 3), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 3									
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	1792 ± 0.0	1792	225 ± 0.0	225	68 ± 0.0	68	8.57	8.57 ± 0.0	8.57
sabre	1172 ± 15.38	1144	342.25 ± 13.38	290	43.56 ± 1.0	43	86.44	5.4 ± 0.05	5.53

name = kagome, size = (5, 5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 0									
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	270 ± 0.0	270	14 ± 0.0	14	164 ± 0.0	164	1.48	1.48 ± 0.0	1.48
sabre	1247 ± 9.78	1247	132.81 ± 3.09	123	128.31 ± 1.72	125	5.12	0.32 ± 0.03	0.37

name = kagome, size = (5, 5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 1									
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	270 ± 0.0	270	14 ± 0.0	14	164 ± 0.0	164	1.47	1.47 ± 0.0	1.47
sabre	338 ± 14.72	345	50.94 ± 3.44	38	110.69 ± 1.34	109	5.47	0.34 ± 0.02	0.3

name = kagome, size = (5, 5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 2									
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	270 ± 0.0	270	14 ± 0.0	14	164 ± 0.0	164	1.47	1.47 ± 0.0	1.47
sabre	364 ± 10.22	350	48.81 ± 4.88	38	110.19 ± 2.16	115	9.7	0.61 ± 0.03	0.56

name = kago	ome, size = (5,	5), circuit_typ	e = quantum_sim	ulation, p =	1, repetitions =	16, optimizatio	on_level = 3			
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	270 ± 0.0	270	14 ± 0.0	14	164 ± 0.0	164	1.47	1.47 ± 0.0	1.47	
sabre	274 ± 9.62	269	37.44 ± 3.25	30	106.25 ± 1.62	108	36.54	2.28 ± 0.03	2.33	
name = kago	ome, size = (5,	5), circuit_typ	e = quantum_sim	mulation, p =	8, repetitions =	16, optimizatio	on_level = 0			
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	2160 ± 0.0	2160	98 ± 0.0	98	164 ± 0.0	164	10.96	10.96 ± 0.0	10.96	
sabre	6068 ± 87.86	5579	557.62 ± 12.5	511	134.88 ± 2.03	134	27.49	1.72 ± 0.06	1.55	

name = kago	me, size = (5, 8	5), circuit_typ	e = quantum_simu	lation, p = 8,	repetitions =	16, optimization_	level = 1		
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	2160 ± 0.0	2160	98 ± 0.0	98	164 ± 0.0	164	11.01	11.01 ± 0.0	11.01
sabre	2124 ± 45.1	2056	341.69 ± 14.66	286	106.5 ± 2.38	118	30.77	1.92 ± 0.05	1.98

name = kagome, size = (5, 5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 2									
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	2160 ± 0.0	2160	98 ± 0.0	98	164 ± 0.0	164	10.93	10.93 ± 0.0	10.93
sabre	1894 ± 32.99	1953	324.31 ± 22.53	265	106.62 ± 2.47	118	44.89	2.81 ± 0.03	2.8

<pre>name = kagome, size = (5, 5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 3</pre>									
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	2160 ± 0.0	2160	98 ± 0.0	98	164 ± 0.0	164	10.88	10.88 ± 0.0	10.88
sabre	1869 ± 23.79	1758	282.88 ± 13.62	236	104.12 ± 1.91	109	158.7	9.92 ± 0.08	9.89

name = kagome, size = (5, 5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 0									
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
	4000 1 0 0	4000	101 1 0 0	404			04.05		04.05
line-graph	4320 ± 0.0	4320	194 ± 0.0	194	164 ± 0.0	164	21.85	21.85 ± 0.0	21.85
sabre	10420 ± 144.88	9561	979 94 ± 18 51	919	133 62 ± 2 94	124	56 99	3.56 ± 0.11	3 42
20016	10420 - 144.00	5501	575.54 - 10.51	515	100.02 = 2.04	124	00.00	5.50 - 0.11	0.42

<pre>name = kagome, size = (5, 5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 1</pre>										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	4320 ± 0.0	4320	194 ± 0.0	194	164 ± 0.0	164	21.73	21.73 ± 0.0	21.73	
sabre	4267 ± 109.67	4443	682.12 ± 23.85	627	107.31 ± 3.03	110	58.8	3.68 ± 0.11	3.58	

name = kagome, size = (5, 5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 2 method av. n_swaps min. n_swap av. depth min. depth av. n_qubits min. qubits total time (s) av. time (s) min. tim	
name = kagome, size = (5, 5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 2 method av. n_swaps min. n_swap av. depth min. depth av. n_qubits min. qubits total time (s) av. time (s) min. tim	
method av. n_swaps min. n_swap av. depth min. depth av. n_qubits min. qubits total time (s) av. time (s) min. tim	
method av. n_swaps min. n_swap av. depth min. depth av. n_qubits min. qubits total time (s) av. time (s) min. tim	
	1e (s)
$11ne-graph 4320 \pm 0.0 4320 194 \pm 0.0 194 164 \pm 0.0 164 21.86 \pm 0.0$	21.86
sabre 3972 ± 57 03 3804 628 44 ± 28 22 526 107 31 ± 2 53 109 85 09 5 32 ± 0 1	5 11
	0.11

name = kagome, size = (5, 5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 3										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	4320 ± 0.0	4320	194 ± 0.0	194	164 ± 0.0	164	21.56	21.56 ± 0.0	21.56	
sabre	3716 ± 30.9	3656	592.5 ± 13.72	562	104.94 ± 1.25	108	296.62	18.54 ± 0.15	18.2	

<pre>name = kagome, size = (7, 7), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 0</pre>											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	498 ± 0.0	498	16 ± 0.0	16	300 ± 0.0	300	2.8	2.8 ± 0.0	2.8		
sabre	3534 ± 31.01	3534	208.06 ± 8.72	175	238.75 ± 3.75	239	10.84	0.68 ± 0.03	0.71		

name = kagome, size = (7, 7), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 1										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	498 ± 0.0	498	16 ± 0.0	16	300 ± 0.0	300	2.86	2.86 ± 0.0	2.86	
sabre	1119 ± 28.5	1092	94.94 ± 7.28	70	202.19 ± 3.62	210	13.88	0.87 ± 0.03	0.8	

name = kagome, size = (7, 7), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 2										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	498 ± 0.0	498	16 ± 0.0	16	300 ± 0.0	300	2.88	2.88 ± 0.0	2.88	
sabre	994 ± 28.53	902	83.25 ± 4.84	68	200.81 ± 3.18	194	30.74	1.92 ± 0.04	1.95	

name = kago	me, size = (7,	7), circuit_typ	e = quantum_si	mulation, p =	1, repetitions =	= 16, optimizati	on_level = 3		
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	498 ± 0.0	498	16 ± 0.0	16	300 ± 0.0	300	2.87	2.87 ± 0.0	2.87
sabre	868 ± 28.66	737	66.25 ± 3.94	53	194.75 ± 2.84	189	135.51	8.47 ± 0.05	8.39

name = kagome, size = (7, 7), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 0										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	3984 ± 0.0	3984	114 ± 0.0	114	300 ± 0.0	300	21.51	21.51 ± 0.0	21.51	
sabre	13056 ± 225.0	13056	766.81 ± 18.31	686	247.75 ± 2.97	251	61.25	3.83 ± 0.12	4.01	

name = kagome, size = (7, 7), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 1										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	3984 ± 0.0	3984	114 ± 0.0	114	300 ± 0.0	300	20.92	20.92 ± 0.0	20.92	
sabre	5532 ± 133.17	4986	569.0 ± 30.89	426	199.12 ± 3.19	207	68.32	4.27 ± 0.11	4.14	

name = kagome, size = (7, 7), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 2										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	3984 ± 0.0	3984	114 ± 0.0	114	300 ± 0.0	300	20.7	20.7 ± 0.0	20.7	
sabre	4758 ± 150.83	4644	511.88 ± 33.91	429	194.75 ± 3.38	199	118.43	7.4 ± 0.07	7.32	

name = kagome, size = (7, 7), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 3											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	3984 ± 0.0	3984	114 ± 0.0	114	300 ± 0.0	300	20.74	20.74 ± 0.0	20.74		
sabre	4702 ± 106.77	4300	460.69 ± 27.72	375	191.88 ± 2.75	200	504.51	31.53 ± 0.21	31.37		

name = kago	me, size = (7, 7), circuit_type	e = quantum_simula	tion, p = 16,	repetitions = 10	<pre>6, optimization_3</pre>	level = 0		
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	7968 ± 0.0	7968	226 ± 0.0	226	300 ± 0.0	300	41.85	41.85 ± 0.0	41.85
sabre	22309 ± 375 19	21499	1238 75 ± 23 39	1182	248 75 ± 2 74	242	116 84	73±019	7 39
20016	22003 - 010.15	21433	1200.10 - 20.00	1102	240.10 - 2.14	242	110.04	1.5 - 0.15	1.00

name = kago	name = kagome, size = (7, 7), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 1											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)			
line-graph	7968 ± 0.0	7968	226 ± 0.0	226	300 ± 0.0	300	41.09	41.09 ± 0.0	41.09			
sabre	9497 ± 345.67	9370	1049.75 ± 55.28	888	194.06 ± 3.81	198	134.06	8.38 ± 0.28	8.61			

name = kagome, size = (7, 7), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 2											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	7968 ± 0.0	7968	226 ± 0.0	226	300 ± 0.0	300	41.59	41.59 ± 0.0	41.59		
0.											
sabre	9742 ± 181.92	9337	978.25 ± 31.91	856	197.69 ± 3.94	206	227.53	14.22 ± 0.41	15.89		

name = kagome, size = (7, 7), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 3											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	7968 ± 0.0	7968	226 ± 0.0	226	300 ± 0.0	300	42.57	42.57 ± 0.0	42.57		
sabre	8452 ± 100.74	8486	899.0 ± 24.78	790	192.25 ± 2.62	200	886.06	55.38 ± 1.31	60.71		

<pre>name = shuriken, size = (1, 1), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 0</pre>											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	8 ± 0.0	8	9 ± 0.0	9	8 ± 0.0	8	0.09	0.09 ± 0.0	0.09		
sabre	8 ± nan	8	8.88 ± 0.19	8	8.0 ± nan	8	0.32	0.02 ± nan	0.02		

name = shuriken, size = (1, 1), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 1											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	8 ± 0.0	8	9 ± 0.0	9	8 ± 0.0	8	0.09	0.09 ± 0.0	0.09		
sabre	6 ± nan	6	7.5 ± 0.25	7	8.0 ± nan	8	0.32	0.02 ± nan	0.02		

name = shuriken, size = (1, 1), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 2											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	8 ± 0.0	8	9 ± 0.0	9	8 ± 0.0	8	0.09	0.09 ± 0.0	0.09		
sabre	6 ± nan	6	7.5 ± 0.25	7	8.0 ± nan	8	0.32	0.02 ± nan	0.02		

<pre>name = shuriken, size = (1, 1), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 3</pre>									
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	8 ± 0.0	8	9 ± 0.0	9	8 ± 0.0	8	0.09	0.09 ± 0.0	0.09
sabre	6 ± nan	6	7.62 ± 0.22	7	8.0 ± nan	8	0.64	0.04 ± nan	0.04

name = shuriken, size = (1, 1), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 0											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	64 ± 0.0	64	65 ± 0.0	65	8 ± 0.0	8	0.74	0.74 ± 0.0	0.74		
sabre	74 ± 1.56	64	67.94 ± 2.31	61	8.0 ± nan	8	1.66	0.1 ± 0.02	0.09		

name = shuriken, size = (1, 1), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 1											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	64 ± 0.0	64	65 ± 0.0	65	8 ± 0.0	8	0.73	0.73 ± 0.0	0.73		
sabre	62 ± nan	62	60.75 ± 0.53	59	8.0 ± nan	8	1.82	0.11 ± 0.02	0.1		

<pre>name = shuriken, size = (1, 1), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 2</pre>											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	64 ± 0.0	64	65 ± 0.0	65	8 ± 0.0	8	0.77	0.77 ± 0.0	0.77		
sabre	62 ± nan	62	60.19 ± 0.69	58	8.0 ± nan	8	2.04	0.13 ± 0.02	0.13		

name = shuriken, size = (1, 1), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 3											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	64 ± 0.0	64	65 ± 0.0	65	8 ± 0.0	8	0.75	0.75 ± 0.0	0.75		
anhra	60 ± non	60	60 60 + 0 66	EQ	9 0 ± non	•	E 7	0.26 + 0.02	0.41		
Sable	02 ± 11ali	02	00.09 - 0.00	56	0.0 ± 11ali	0	5.7	0.30 - 0.03	0.41		

<pre>name = shuriken, size = (1, 1), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 0</pre>												
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)			
line-graph	128 ± 0.0	128	129 ± 0.0	129	8 ± 0.0	8	1.5	1.5 ± 0.0	1.5			
sabre	138 ± 2.94	138	140.56 ± 3.66	128	8.0 ± nan	8	3.42	0.21 ± 0.03	0.18			

method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	128 ± 0.0	128	129 ± 0.0	129	8 ± 0.0		1.53	1.53 ± 0.0	1.53
sabre	126 ± nan	126	121.12 ± 1.06	119	8.0 ± nan	8	3.57	0.22 ± 0.02	0.2
name = shur	iken, size = (1	, 1), circuit_t		imulation, p =	16, repetition	s = 16, optimizat	tion_level = 2		
nethod	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
ine-graph	128 ± 0.0	128	129 ± 0.0	129	8 ± 0.0	8	1.58	1.58 ± 0.0	1.58
line-graph sabre name = shur method	128 ± 0.0 126 ± nan 	128 126 	129 ± 0.0 121.06 ± 0.88 	129 118 	<pre>8 ± 0.0 8.0 ± nan 4.0 5.0 5.0 5.0 5.0 5.0 5.0 5.0 5.0 5.0 5</pre>	8 8 9 = 16, optimiza min. qubits	1.58 4.24 	1.58 ± 0.0 0.26 ± 0.03	1.5 0.2 min. time (s
line-graph sabre name = shur method line-graph	128 ± 0.0 126 ± nan tiken, size = (1 av. n_swaps 128 ± 0.0	128 126	129 ± 0.0 121.06 ± 0.88 	129 118 	8 ± 0.0 8.0 ± nan • 16, repetition av. n_qubits 	8 8 5 = 16, optimizat min. qubits 	1.58 4.24 tion_level = 3 total time (s) 1.5	1.58 ± 0.0 0.26 ± 0.03 av. time (s)	1.5 0.2 min. time (s
line-graph sabre name = shur method line-graph sabre	128 ± 0.0 126 ± nan 126 ± nan 126 ± nan 128 ± 0.0 128 ± 0.0 126 ± nan	128 126 	129 ± 0.0 121.06 ± 0.88 	129 118 imulation, p = min. depth 129 116	8 ± 0.0 8.0 ± nan • 16, repetition av. n_qubits 	8 8 5 = 16, optimiza min. qubits 	1.58 4.24 tion_level = 3 total time (s) 	1.58 ± 0.0 0.26 ± 0.03 av. time (s) 	1.54 0.24 min. time (s)
line-graph sabre 	128 ± 0.0 126 ± nan 126 ± nan 128 ± 0.0 128 ± 0.0 128 ± 0.0 126 ± nan	128 126 ., 1), circuit_t min. n_swap 	129 ± 0.0 121.06 ± 0.88 	129 118 imulation, p = min. depth 	<pre>8 # 0.0 8.0 # nan 16, repetition av. n_qubits 6 # 0.0 8.0 # nan 1. repetitions 1. repetitio</pre>	8 8 5 = 16, optimizat min. qubits 	1.58 4.24 tion_level = 3 total time (s) 1.5 11.06	1.58 ± 0.0 0.26 ± 0.03 av. time (s) 1.5 ± 0.0 0.69 ± 0.02	1.53 0.24 min. time (s) 1.5 0.73
line-graph sabre 	128 ± 0.0 126 ± nan 126 ± nan 128 ± 0.0 128 ± 0.0 128 ± 0.0 126 ± nan 126 ± nan 126 ± nan	128 126 ., 1), circuit_t min. n_swap 	129 ± 0.0 121.06 ± 0.88 	129 118 imulation, p = min. depth 129 116 imulation, p = min. depth	<pre>8 # 0.0 8.0 # nan * 16, repetition av. n_qubits </pre>	8 8 5 = 16, optimizat min. qubits 8 8 8 = 16, optimizat: min. qubits	1.58 4.24 tion_level = 3 total time (s) 1.5 11.06 total time (s)	1.58 ± 0.0 0.26 ± 0.03 av. time (s) 	1.53 0.24 min. time (s) 1.5 0.7 min. time (s)
line-graph sabre name = shur nethod line-graph sabre name = shur nethod 	128 ± 0.0 126 ± nan 126 ± nan 126 ± nan 128 ± 0.0 128 ± 0.0 128 ± 0.0 126 ± nan 128 ± 0.0 126 ± nan 128 ± 0.0 126 ± nan 128 ± 0.0 126 ± nan	128 126 	129 ± 0.0 121.06 ± 0.88 	129 118 	<pre>8 # 0.0 8.0 # nan • 16, repetition av. n_qubits </pre>	8 8 5 = 16, optimizat min. qubits 8 8 9 = 16, optimizat: min. qubits 	1.58 4.24 tion_level = 3 total time (s) 	1.58 ± 0.0 0.26 ± 0.03 av. time (s) 1.5 ± 0.0 0.69 ± 0.02 av. time (s) 1.0 ± 0.0	1.53 0.24 min. time (s) 1.5 0.7 min. time (s)

<pre>name = shuriken, size = (3, 3), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 1</pre>										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	130 ± 0.0	130	12 ± 0.0	12	84 ± 0.0	84	0.89	0.89 ± 0.0	0.89	
sabre	177 ± 6.53	187	34.5 ± 2.94	27	64.88 ± 1.19	67	2.79	0.17 ± 0.02	0.16	

name = shuriken, size = (3, 3), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 2											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	130 ± 0.0	130	12 ± 0.0	12	84 ± 0.0	84	0.89	0.89 ± 0.0	0.89		
sabre	169 ± 7.66	123	34.5 ± 2.44	26	64.5 ± 1.0	63	4.43	0.28 ± 0.02	0.24		

nethod	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
Line-graph	130 ± 0.0	130	12 ± 0.0	12	84 ± 0.0	84	1	1.0 ± 0.0	1
abre	138 ± 6.41	94	29.44 ± 2.53	20	62.81 ± 0.84	63	14.29	0.89 ± 0.03	0.94
ame = shur	riken, size = (3	, 3), circuit_t	ype = quantum_s:	imulation, p	= 8, repetitions	= 16, optimizat	ion_level = 0		
ethod	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s
ine-graph	1040 ± 0.0	1040	89 ± 0.0	89	84 ± 0.0	84	7.06	7.06 ± 0.0	7.0
abre	2778 ± 36.62	2575	400.94 ± 8.12	362	83.81 ± 0.19	84	18.11	1.13 ± 0.03	1.1
ame = shur ethod	riken, size = (3 av. n_swaps	, 3), circuit_t min. n_swap	ype = quantum_si av. depth	imulation, p min. dept	= 8, repetitions h av. n_qubits	= 16, optimizat	ion_level = 1 total time (s.) av. time (s)	min. time
ame = shur wethod ine-graph abre	<pre>iken, size = (3 av. n_swaps</pre>	, 3), circuit_t min. n_swap 1040 1039	<pre>ype = quantum_s: av. depth </pre>	imulation, p min. dept 	 = 8, repetitions h av. n_qubits 	= 16, optimizat min. qubits 	ion_level = 1 total time (s 6.9 16.0) av. time (s) 	min. time (6. 0.
ame = shur ethod ine-graph abre	iken, size = (3 av. n_swaps 	, 3), circuit_t min. n_swap 1040 1039	<pre>ype = quantum_s: av. depth </pre>	imulation, p min. dept 	<pre>= 8, repetitions h av. n_qubits 9 84 * 0.0 5 64.94 * 1.41</pre>	= 16, optimizat. min. qubits 	ion_level = 1 total time (s) av. time (s) 2 6.92 ± 0.0 5 1.0 ± 0.04	min. time (6. 0.
ame = shur ethod ine-graph abre ame = shur ethod	<pre>iken, size = (3 av. n_swaps 1040 ± 0.0 1268 ± 32.28 iken, size = (3 av. n_swaps</pre>	<pre>, 3), circuit_t min. n_swap</pre>	<pre>ype = quantum_s: av. depth </pre>	imulation, p min. dept 8 20 	<pre>= 8, repetitions h av. n_qubits</pre>	<pre>= 16, optimizat min. qubits 84 63 = 16, optimizat min. qubits</pre>	<pre>ion_level = 1 total time (s</pre>) av. time (s) 2 6.92 ± 0.0 5 1.0 ± 0.04	min. time 0 6 0 min. time 1
ame = shur ethod 	<pre>iken, size = (3 av. n_swaps</pre>	<pre>, 3), circuit_t min. n_swap</pre>	<pre>ype = quantum_s: av. depth </pre>	imulation, p min. dept 8 20 	<pre>= 8, repetitions h av. n_qubits</pre>	 = 16, optimizat. min. qubits 84 63 = 16, optimizat. min. qubits 	<pre>ion_level = 1 total time (s</pre>) av. time (s) 2 6.92 ± 0.0 5 1.0 ± 0.04) av. time (s)	min. time (6 0 min. time (
ame = shur sthod 	<pre>iken, size = (3 av. n_swaps 1040 ± 0.0 1268 ± 32.28 iken, size = (3 av. n_swaps 1040 ± 0.0</pre>	<pre>, 3), circuit_t min. n_swap</pre>	<pre>ype = quantum_s: av. depth </pre>	imulation, p min. dept 8 20 20 imulation, p min. dept	<pre>= 8, repetitions h av. n_qubits</pre>	 = 16, optimizat. min. qubits 84 63 = 16, optimizat. min. qubits 84 	<pre>ion_level = 1 total time (s</pre>	<pre>) av. time (s) </pre>	min. time (6 0 min. time (
me = shun ne-graph .bre 	<pre>iken, size = (3 av. n_swaps</pre>	<pre>, 3), circuit_t min. n_swap</pre>	<pre>ype = quantum_s: av. depth </pre>	imulation, p min. dept 20 	<pre>= 8, repetitions h av. n_qubits</pre>	 = 16, optimizat. min. qubits 84 63 = 16, optimizat. min. qubits 84 65 	<pre>ion_level = 1 total time (s</pre>	<pre>) av. time (s) </pre>	min. time 6 0
ame = shur ethod abre ame = shur ethod ine-graph abre	<pre>iken, size = (3 av. n_swaps</pre>	<pre>, 3), circuit_t min. n_swap 1040 1039 , 3), circuit_t min. n_swap 1040 926</pre>	<pre>ype = quantum_s: av. depth </pre>	imulation, p min. dept 20 imulation, p min. dept 317	<pre>= 8, repetitions h av. n_qubits</pre>	 = 16, optimizat. min. qubits 84 63 = 16, optimizat. min. qubits 84 65 	<pre>ion_level = 1 total time (s</pre>	<pre>) av. time (s) 2 6.92 ± 0.0 5 1.0 ± 0.04) av. time (s) 2 6.92 ± 0.0 1.41 ± 0.05</pre>	min. time 6 0 min. time 6 1
ame = shur sthod 	<pre>iken, size = (3 av. n_swaps</pre>	<pre>, 3), circuit_t min. n_swap</pre>	<pre>ype = quantum_s: av. depth </pre>	imulation, p min. dept 8 20 	<pre>= 8, repetitions h av. n_qubits</pre>	<pre>= 16, optimizat: min. qubits</pre>	<pre>ion_level = 1 total time (s</pre>	<pre>) av. time (s) </pre>	min. time 6 0

 line-graph
 1040 ± 0.0
 1040
 89 ± 0.0
 89
 84 ± 0.0
 84
 7.06
 7.06 ± 0.0
 7.06

 sabre
 966 ± 14.37
 892
 192.5 ± 13.59
 145
 64.44 ± 1.09
 68
 79.12
 4.94 ± 0.07
 5.05

<pre>name = shuriken, size = (3, 3), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 0</pre>												
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)			
line-graph	2080 ± 0.0	2080	177 ± 0.0	177	84 ± 0.0	84	13.7	13.7 ± 0.0	13.7			
sabre	5067 ± 112.35	4490	774.94 ± 15.84	706	83.94 ± 0.16	84	31.53	1.97 ± 0.08	2.1			

name = shuriken, size = (3, 3), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 1										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	2080 ± 0.0	2080	177 ± 0.0	177	84 ± 0.0	84	14.13	14.13 ± 0.0	14.13	
sabre	2035 ± 50.28	1861	450.12 ± 38.78	302	65.19 ± 1.28	66	35.28	2.2 ± 0.06	2.19	

<pre>name = shuriken, size = (3, 3), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 2</pre>											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	2080 ± 0.0	2080	177 ± 0.0	177	84 ± 0.0	84	13.71	13.71 ± 0.0	13.71		
sabre	2080 ± 41.66	1890	398.25 ± 19.59	334	65.38 ± 0.94	63	44.87	2.8 ± 0.06	2.84		

		2)			10	- 10			
name = snur.	iken, size = (3,	, S), CIFCUIT_t	.ype = quantum_si	mulation, p =	16, repetitions	= 16, optimizati	.on_ievei = 3		
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line ment	2020 + 0.0	2080	177 + 0.0	177	84 + 0.0		12.00	12 20 + 0 0	12.00
iine-graph	2080 ± 0.0	2080	177 ± 0.0	1//	84 ± 0.0	04	13.69	13.69 ± 0.0	13.69
sabre	1826 ± 35.68	1703	373.75 ± 26.89	303	64.44 ± 1.53	61	150.75	9.42 ± 0.09	9.59

name = shuriken, size = (5, 5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 0											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	418 ± 0.0	418	14 ± 0.0	14	240 ± 0.0	240	2.76	2.76 ± 0.0	2.76		
sabre	1722 ± 13.92	1706	117.56 ± 4.03	100	237.69 ± 0.81	239	7.59	0.47 ± 0.03	0.43		

name = shuriken, size = (5, 5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 1											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	418 ± 0.0	418	14 ± 0.0	14	240 ± 0.0	240	2.61	2.61 ± 0.0	2.61		
sabre	993 ± 25.43	918	85.12 ± 4.62	70	179.31 ± 2.32	182	11.16	0.7 ± 0.03	0.8		

name = shuriken, size = (5, 5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 2											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	418 ± 0.0	418	14 ± 0.0	14	240 ± 0.0	240	2.58	2.58 ± 0.0	2.58		
sabre	940 ± 25.45	830	83.06 ± 7.81	54	179.06 ± 3.62	194	25.23	1.58 ± 0.02	1.55		

<pre>name = shuriken, size = (5, 5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 3</pre>										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	418 ± 0.0	418	14 ± 0.0	14	240 ± 0.0	240	2.6	2.6 ± 0.0	2.6	
sabre	702 ± 15.83	712	66.81 ± 6.6	45	176.12 ± 2.69	183	105.59	6.6 ± 0.04	6.58	

name = shur	name = shuriken, size = (5, 5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 0											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)			
line-graph	3344 ± 0.0	3344	105 ± 0.0	105	240 ± 0.0	240	19.21	19.21 ± 0.0	19.21			
cabro	10085 + 62 18	10397	654 81 + 12 53	619	238 31 + 0 59	238	53 61	3 35 + 0 12	3 66			
Sabre	10000 - 02.10	10001	004.01 - 12.00	015	200.01 - 0.00	200	55.01	5.55 - 0.12	5.00			

name = shuriken, size = (5, 5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 1											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	3344 ± 0.0	3344	105 ± 0.0	105	240 ± 0.0	240	19.42	19.42 ± 0.0	19.42		
sabre	4949 ± 172.47	4147	581.88 ± 36.5	461	178.88 ± 2.91	173	60.32	3.77 ± 0.09	3.72		

name = shuriken, size = (5, 5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 2											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	3344 ± 0.0	3344	105 ± 0.0	105	240 ± 0.0	240	19.65	19.65 ± 0.0	19.65		
sabre	4321 ± 153.74	3792	566.06 ± 43.95	414	173.81 ± 2.09	171	109.02	6.81 ± 0.22	7.04		

name = shuriken, size = (5, 5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 3											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	3344 ± 0.0	3344	105 ± 0.0	105	240 ± 0.0	240	19.33	19.33 ± 0.0	19.33		
sabre	3405 ± 65.66	3312	408.12 ± 21.62	343	170.62 ± 2.41	165	429.48	26.84 ± 0.98	25.93		

name = shuriken, size = (5, 5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 0											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	6688 ± 0.0	6688	209 ± 0.0	209	240 ± 0.0	240	38.44	38.44 ± 0.0	38.44		
0 1											
sabre	18976 ± 105.01	18874	1155.94 ± 14.25	1095	238.88 ± 0.38	238	108.03	6.75 ± 0.22	6.3		

name = shuriken, size = (5, 5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 1										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	6688 ± 0.0	6688	209 ± 0.0	209	240 ± 0.0	240	38.54	38.54 ± 0.0	38.54	
sabre	8889 ± 336.48	6839	1004.38 ± 82.18	832	174.88 ± 2.72	176	111.6	6.98 ± 0.17	6.8	

name = shuriken, size = (5, 5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 2											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line_graph	6688 + 0 0	6688	209 + 0 0	209	240 + 0 0	240	38 42	38 42 + 0 0	38.42		
1116-Brabu	0000 - 0.0	0088	209 - 0.0	209	240 - 0.0	240	30.42	38.42 - 0.0	30.42		
sabre	6943 ± 155.17	6914	836.81 ± 46.97	673	173.06 ± 1.69	173	183.47	11.47 ± 0.2	12.22		

<pre>name = shuriken, size = (5, 5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 3</pre>											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	6688 ± 0.0	6688	209 ± 0.0	209	240 ± 0.0	240	38.14	38.14 ± 0.0	38.14		
sabre	6644 ± 94.12	6326	667.75 ± 19.72	599	170.81 ± 2.19	165	693.94	43.37 ± 1.15	43.48		

name = shuriken, size = (7, 7), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 0										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	850 ± 0.0	850	14 ± 0.0	14	476 ± 0.0	476	5.19	5.19 ± 0.0	5.19	
sabre	4169 ± 24.72	4173	143.25 ± 4.62	127	471.81 ± 1.22	474	16.26	1.02 ± 0.05	1.13	

name = shuriken, size = (7, 7), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 1											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	850 ± 0.0	850	14 ± 0.0	14	476 ± 0.0	476	5.2	5.2 ± 0.0	5.2		
sabre	2783 ± 64.06	2792	196.44 ± 13.5	154	343.5 ± 3.28	349	48.94	3.06 ± 0.24	2.78		

name = shuriken, size = (7, 7), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 2										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	850 ± 0.0	850	14 ± 0.0	14	476 ± 0.0	476	5.31	5.31 ± 0.0	5.31	
sabre	2665 ± 65.58	2535	170.75 ± 14.04	127	350.5 ± 5.34	352	132.45	8.28 ± 0.26	7.5	

name = shuriken, size = (7, 7), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 3									
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	850 ± 0.0	850	14 ± 0.0	14	476 ± 0.0	476	5.22	5.22 ± 0.0	5.22
sabre	2079 ± 57.9	1891	124.12 ± 8.31	93	342.06 ± 3.19	354	628.32	39.27 ± 0.78	37.63

name = shuwikan sime = $(7, 7)$ simulit ture = supplum simulation $n = 0$ variations = 16 estimization layel = 0										
none = chunikan size = $(7, 7)$ sincuit ture = quantum simulation $n = 9$ repetitions = 16 entimization level = 0										
name = shuriken, size = (7, 7), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 0										
method av. n_swaps min. n_swap av. depth min. depth av. n_qubits min. qubits total time (s) av. time (s) min. t	ime (s)									
line-graph 6800 * 0.0 6800 105 * 0.0 105 476 * 0.0 476 38.58 * 0.0	38.58									
sabre 25487 * 146.92 25288 875.19 * 21.34 785 474.06 * 1.12 475 111.3 6.96 * 0.17	6.82									

name = shuriken, size = (7, 7), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 1										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	6800 ± 0.0	6800	105 ± 0.0	105	476 ± 0.0	476	38.22	38.22 ± 0.0	38.22	
sabre	14038 ± 501.85	11572	1110.0 ± 72.38	836	349.44 ± 5.69	345	167.86	10.49 ± 0.26	10.87	

name = shuriken, size = (7, 7), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 2										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	6800 ± 0.0	6800	105 ± 0.0	105	476 ± 0.0	476	37.94	37.94 ± 0.0	37.94	
sabre	11889 ± 326.3	12192	1092.25 ± 51.37	922	343.62 ± 5.27	347	352.13	22.01 ± 0.34	20.89	

		-								
name = shuriken, size = (7, 7), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 3										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	6800 ± 0.0	6800	105 ± 0.0	105	476 ± 0.0	476	38.25	38.25 ± 0.0	38.25	
sabre	9431 ± 295.31	9200	791.69 ± 50.65	589	336.75 ± 4.19	333	1612.88	100.8 ± 1.84	99.38	

name = shur	iken, size = (7,	7), circuit_ty	pe = quantum_simu	lation, p = 16	6, repetitions =	16, optimization	_level = 0		
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line graph	12600 + 0 0	12600	200 + 0 0	200	476 + 0.0	476	75 20	75 28 ± 0 0	75 20
11ne-graph	13600 ± 0.0	13600	209 ± 0.0	209	476 ± 0.0	470	10.30	75.36 ± 0.0	/5.30
sabre	45512 ± 312.53	44697	1457.12 ± 33.54	1343	473.31 ± 1.06	474	207.14	12.95 ± 0.2	13.88

name = shur	riken, size = (7,	7), circuit_t	ype = quantum_simul	lation, p = 16	, repetitions =	16, optimization	level = 1		
method	av. n_swaps	min. n_swa	up av. depth	min. depth	n av. n_qubits	min. qubits	total time (s) av. time (s)	min. time (s)
line-graph	13600 ± 0.0	1360	0 209 ± 0.0	209	9 476 ± 0.0	476	75.8	2 75.82 ± 0.0	75.82
sabre	24677 ± 1008.11	1921	3 1943.62 ± 179.76	5 1252	2 342.62 ± 4.44	334	295.3	1 18.46 ± 0.33	19.18
name = shur	riken, size = (7,	7), circuit_t	ype = quantum_simul	Lation, p = 16	, repetitions =	16, optimization	level = 2		
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	13600 ± 0.0	13600	209 ± 0.0	209	476 ± 0.0	476	75.21	75.21 ± 0.0	75.21
sabre	22096 ± 699.86	19115	1676.12 ± 96.85	1293	348.44 ± 4.94	347	573.52	35.84 ± 0.93	38.88
name = shur	riken, size = (7,	7), circuit_t	ype = quantum_simul	lation, p = 16	, repetitions =	16, optimization	level = 3		
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	13600 ± 0.0	13600	209 ± 0.0	209	476 ± 0.0	476	76.74	76.74 ± 0.0	76.74
sabre	18148 ± 348.43	17064	1369.19 ± 66.21	1099	336.31 ± 4.62	323	2519.55	157.47 ± 3.75	146.54

7.0.2 Quantum simulation, checkerboard, agianst SABRE

```
import line_graph_routing as lgr # Loading these makes these cells stand-alone
import pickle
settings=[]
for name in ['checkerboard']:
   for side in [i+0.5 for i in range(1,9,2)]:
        for p in [1,8,16]:
            for optimization_level in range(4):
                setting={'name':name,
                     'size': (side,side),
                     'circuit_type': 'quantum_simulation',
                     'p': p,
                     'repetitions' : 16,
                     'optimization_level' : optimization_level,
                     'methods' : ['sabre']
                    }
                settings.append(setting)
## Uncomment to rerun benchmarks. This takes a couple of hours.
#results=[]
#for setting in settings:
    result=lqr.benchmark(**setting)
#
```

```
# results.append(result)
# lgr.print_benchmark(result)
#
#with open('benchmark_data/checkerboard.pkl','wb') as f:
# pickle.dump(results,f)
#Load previously obtained results from disk and show them.
import pickle
with open('benchmark_data/checkerboard.pkl','rb') as f:
results=pickle.load(f)
for result in results:
    lgr.print_benchmark(result)
```

name = chec	ckerboard, size =	(1.5, 1.5), o	ircuit_type =	quantum_simula	ation, p = 1, re	petitions = 16,	optimization_leve	1 = 0	
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	42 ± 0.0	42	27 ± 0.0	27	21 ± 0.0	21	0.17	0.17 ± 0.0	0.17
sabre	45 ± 1.28	47	32.44 ± 1.97	27	19.06 ± 0.66	21	0.99	0.06 ± 0.03	0.03
name = chec	ckerboard, size =	= (1.5, 1.5), c	ircuit_type =	quantum_simula	ation, p = 1, re	petitions = 16,	optimization_leve	1 = 1	
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	42 ± 0.0	42	27 ± 0.0	27	21 ± 0.0	21	0.37	0.37 ± 0.0	0.37
sabre	32 ± 1.44	22	26.56 ± 2.0	19	17.69 ± 0.66	16	0.67	0.04 ± 0.03	0.03
name = chec	ckerboard, size =	= (1.5, 1.5), c	ircuit_type =	quantum_simula	ation, p = 1, re	petitions = 16,	optimization_leve	1 = 2	
method	av. n_swaps	min. n_swap	av. deptn	min. deptn	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	42 ± 0.0	42	27 ± 0.0	27	21 ± 0.0	21	0.35	0.35 ± 0.0	0.35
sabre	28 ± 0.84	31	23.81 ± 1.81	19	17.94 ± 0.87	21	1.34	0.08 ± 0.04	0.04
name = chec	ckerboard, size =	= (1.5, 1.5), c	ircuit_type =	quantum_simula	ation, p = 1, re	petitions = 16,	optimization_leve	1 = 3	
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	42 ± 0.0	42	27 ± 0.0	27	21 ± 0.0	21	0.17	0.17 ± 0.0	0.17
sabre	28 ± 0.78	25	25.31 ± 2.03	18	16.81 ± 0.41	18	2.56	0.16 ± 0.04	0.11

name = checkerboard, size = (1.5, 1.5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 0

				-			-		
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	322 ± 0.0	322	188 ± 0.0	188	21 ± 0.0	21	1.89	1.89 ± 0.0	1.89
sabre	362 ± 7.38	343	235.0 ± 6.03	214	21.0 ± nan	21	4.87	0.3 ± 0.06	0.18

name = checkerboard, size = (1.5, 1.5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 1										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	322 ± 0.0	322	188 ± 0.0	188	21 ± 0.0	21	1.98	1.98 ± 0.0	1.98	
TING Brahn	022 - 010	022	100 - 010	100	21 - 010		1100	1100 - 010	1.00	
sabre	239 ± 3.71	240	197.25 ± 7.75	175	17.94 ± 0.84	16	4.99	0.31 ± 0.06	0.21	

name = checkerboard, size = (1.5, 1.5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 2										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	322 ± 0.0	322	188 ± 0.0	188	21 ± 0.0	21	1.94	1.94 ± 0.0	1.94	
sabre	237 ± 2.75	227	183.5 ± 7.31	160	17.12 ± 0.5	16	5.84	0.36 ± 0.05	0.28	

name = checkerboard, size = (1.5, 1.5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 3										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	322 ± 0.0	322	188 ± 0.0	188	21 ± 0.0	21	1.95	1.95 ± 0.0	1.95	
sabre	217 ± 2.41	221	187.5 ± 4.34	175	16.62 ± 0.44	16	17.58	1.1 ± 0.01	1.08	

name = checkerboard, size = (1.5, 1.5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 0											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	642 ± 0.0	642	372 ± 0.0	372	21 ± 0.0	21	3.85	3.85 ± 0.0	3.85		
sabre	763 ± 20.84	657	485.38 ± 12.94	417	21.0 ± nan	21	8.07	0.5 ± 0.05	0.51		

name = checkerboard, size = (1.5, 1.5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 1										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	642 ± 0.0	642	372 ± 0.0	372	21 ± 0.0	21	3.18	3.18 ± 0.0	3.18	
sabre	471 ± 9.91	469	379.56 ± 17.02	311	18.62 ± 0.88	21	7.55	0.47 ± 0.04	0.4	

name = checkerboard, size = (1.5, 1.5), circuit_type = quantum_simulation, p = 1	6, repetitions = 16, optimization_level = 2

method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	642 ± 0.0	642	372 ± 0.0	372	21 ± 0.0	21	3.22	3.22 ± 0.0	3.22
sabre	463 ± 5.31	466	376.0 ± 8.98	346	17.5 ± 0.81	21	10.29	0.64 ± 0.04	0.52

name = checkerboard, size = (1.5, 1.5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 3											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	642 ± 0.0	642	372 ± 0.0	372	21 ± 0.0	21	3.23	3.23 ± 0.0	3.23		
sabre	454 ± 2.81	447	353.69 ± 12.16	321	16.88 ± 0.5	16	29.5	1.84 ± 0.03	1.96		

name = checkerboard, size = (3.5, 3.5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 0											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	271 ± 0.0	271	30 ± 0.0	30	97 ± 0.0	97	0.92	0.92 ± 0.0	0.92		
sabre	550 ± 9.03	548	104.56 ± 3.78	90	88.5 ± 1.56	92	2.56	0.16 ± 0.02	0.14		

name = checkerboard, size = (3.5, 3.5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 1										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	271 ± 0.0	271	30 ± 0.0	30	97 ± 0.0	97	1.01	1.01 ± 0.0	1.01	
sabre	257 ± 7.69	233	63.81 ± 4.25	48	69.0 ± 1.38	68	3.03	0.19 ± 0.02	0.17	

name = checkerboard, size = (3.5, 3.5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 2										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	271 ± 0.0	271	30 ± 0.0	30	97 ± 0.0	97	1.01	1.01 ± 0.0	1.01	
sabre	228 ± 6.72	234	57.94 ± 3.75	48	70.56 ± 1.25	71	4.87	0.3 ± 0.03	0.26	

name = checkerboard, size = (3.5, 3.5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 3											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
1	071 + 0 0	071	20 + 0 0	20	07 + 0 0	07	1 04	1 04 + 0 0	1 04		
11ne-graph	2/1 ± 0.0	2/1	30 ± 0.0	30	97 ± 0.0	97	1.04	1.04 ± 0.0	1.04		
sabre	180 ± 2 97	187	45 44 ± 2 72	37	67 38 ± 0 97	70	16 94	1.06 ± 0.03	1.1		
Dubro	100 - 2101	101	10111 - 2112	0.	01100 - 0101		10101	1100 - 0100			

name = chec	name = checkerboard, size = (3.5, 3.5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 0												
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)				
line-graph	2077 ± 0.0	2077	198 ± 0.0	198	97 ± 0.0	97	7.69	7.69 ± 0.0	7.69				
sabre	3416 ± 49.08	3159	597.62 ± 15.12	544	92.62 ± 0.75	92	18.07	1.13 ± 0.05	1.19				

name = check	ame = checkerboard, size = (3.5, 3.5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 1											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)			
line-graph	2077 ± 0.0	2077	198 ± 0.0	198	97 ± 0.0	97	7.68	7.68 ± 0.0	7.68			
sabre	1870 ± 26.84	1894	481.25 ± 11.75	444	70.19 ± 1.84	68	20.33	1.27 ± 0.05	1.33			

name = checkerboard, size = (3.5, 3.5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 2											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	2077 ± 0.0	2077	198 ± 0.0	198	97 ± 0.0	97	7.91	7.91 ± 0.0	7.91		
sabre	1785 ± 15.44	1756	456.5 ± 14.31	404	69.56 ± 1.16	73	29.02	1.81 ± 0.05	1.86		

nome = chec	horboard size =	(2 5 2 5)	inquit turno = qu	ontum gimulati	on n = 9 ronot	titions = 16 ont	imigation loval =	2				
name = cnec.												
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)			
line-graph	2077 ± 0.0	2077	198 ± 0.0	198	97 ± 0.0	97	7.73	7.73 ± 0.0	7.73			
sabre	1740 ± 16.03	1731	446.75 ± 14.74	387	69.12 ± 2.0	76	104.9	6.56 ± 0.05	6.48			

name = chec	name = checkerboard, size = (3.5, 3.5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 0											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)			
line-graph	4141 ± 0.0	4141	390 ± 0.0	390	97 ± 0.0	97	14.99	14.99 ± 0.0	14.99			
sabre	6142 ± 72.56	5875	1117.75 ± 21.0	1021	94.0 ± nan	94	36.66	2.29 ± 0.13	2.24			

name = chec	name = checkerboard, size = (3.5, 3.5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 1											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)			
line-graph	4141 ± 0.0	4141	390 ± 0.0	390	97 ± 0.0	97	15.05	15.05 ± 0.0	15.05			
sabre	3639 ± 29.19	3753	996.19 ± 20.97	934	70.0 ± 1.38	67	41.08	2.57 ± 0.09	2.48			

name = chec	kerboard, size	= (3.5, 3.5),	circuit_type = q	uantum_simulat	ion, p = 16, rep	petitions = 16,	optimization_level	. = 2	
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	4141 ± 0.0	4141	390 ± 0.0	390	97 ± 0.0	97	15.66	15.66 ± 0.0	15.66
sabre	3684 ± 25.79	3608	949.5 ± 19.69	893	70.06 ± 2.69	71	57.58	3.6 ± 0.1	3.48

name = chec	me = checkerboard, size = (3.5, 3.5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 3										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	4141 ± 0.0	4141	390 ± 0.0	390	97 ± 0.0	97	15.64	15.64 ± 0.0	15.64		
sabre	3580 ± 20.27	3621	924.56 ± 17.07	854	67.0 ± 1.47	65	207.53	12.97 ± 0.12	12.8		

name = checkerboard, size = (5.5, 5.5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 0											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	682 ± 0.0	682	34 ± 0.0	34	221 ± 0.0	221	2.48	2.48 ± 0.0	2.48		
sabre	1770 ± 17.16	1791	138.56 ± 4.66	121	211.31 ± 1.47	213	6.81	0.43 ± 0.04	0.5		

name = chec	ame = checkerboard, size = (5.5, 5.5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 1												
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)				
line-graph	682 ± 0.0	682	34 ± 0.0	34	221 ± 0.0	221	2.45	2.45 ± 0.0	2.45				
sabre	678 ± 30.68	678	107.06 ± 7.62	77	160.38 ± 1.34	160	9.63	0.6 ± 0.04	0.65				

name = chec	ame = checkerboard, size = (5.5, 5.5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 2											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)			
line-graph	682 ± 0.0	682	34 ± 0.0	34	221 ± 0.0	221	2.47	2.47 ± 0.0	2.47			
sabre	808 ± 21.59	773	102.81 ± 7.0	82	159.62 ± 2.76	160	18.67	1.17 ± 0.03	1.07			

name = checkerboard, size = (5.5, 5.5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 3											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	682 ± 0.0	682	34 ± 0.0	34	221 ± 0.0	221	2.41	2.41 ± 0.0	2.41		
sabre	652 ± 19.28	608	83.75 ± 6.44	58	156.94 ± 1.84	158	80.01	5.0 ± 0.04	5.14		

$name = checkerboard, arze = (0.0, 0.0), cricarc_cype = quancum_armutation, p = 0, repetitiona = 10, optimization_rever =$	name =	= checkerboard,	size = (5	5.5, 5.5),	circuit_type =	quantum_simulation,	p = 8,	repetitions = 16	, optimization_level = (
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method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	5148 ± 0.0	5148	223 ± 0.0	223	221 ± 0.0	221	18.16	18.16 ± 0.0	18.16
sabre	10367 ± 118.27	10404	823.12 ± 18.63	746	217.38 ± 0.62	219	46.53	2.91 ± 0.16	2.8

name = checkerboard, size = (5.5, 5.5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 1										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	5148 ± 0.0	5148	223 ± 0.0	223	221 ± 0.0	221	18.72	18.72 ± 0.0	18.72	
sabre	5644 ± 120.68	5373	800.31 ± 38.98	667	159.44 ± 2.66	166	52.7	3.29 ± 0.1	3.38	

name = checkerboard, size = (5.5, 5.5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 2										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	5148 ± 0.0	5148	223 ± 0.0	223	221 ± 0.0	221	18.19	18.19 ± 0.0	18.19	
sabre	5241 ± 43.19	5091	734.25 ± 22.25	666	157.81 ± 3.06	171	89.44	5.59 ± 0.12	5.86	

name = checkerboard, size = (5.5, 5.5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 3										
method	av. n_swaps	min. n_s	swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	5148 ± 0.0	1	5148	223 ± 0.0	223	221 ± 0.0	221	18.3	18.3 ± 0.0	18.3
sabre	5054 ± 30.28	4	4839	723.25 ± 16.98	656	157.69 ± 2.69	159	363.83	22.74 ± 0.2	23.19

name = checkerboard, size = (5.5, 5.5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 0										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	10252 ± 0.0	10252	439 ± 0.0	439	221 ± 0.0	221	36.56	36.56 ± 0.0	36.56	
sabre	18950 ± 206.58	18281	1546.5 ± 24.0	1464	217.38 ± 0.78	218	98.11	6.13 ± 0.21	6.55	

name = checkerboard, size = (5.5, 5.5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 1										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	10252 ± 0.0	10252	439 ± 0.0	439	221 ± 0.0	221	36.48	36.48 ± 0.0	36.48	
sabre	10695 ± 85.99	10470	1609.69 ± 39.03	1439	161.44 ± 2.88	158	113.03	7.06 ± 0.21	6.69	

name = chec	name = checkerboard, size = (5.5, 5.5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 2										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	10252 ± 0.0	10252	439 ± 0.0	439	221 ± 0.0	221	36.69	36.69 ± 0.0	36.69		
sabre	10265 ± 67.73	10411	1565.5 ± 31.69	1469	162.69 ± 4.19	163	171.93	10.75 ± 0.22	10.26		

name = checkerboard, size = (5.5, 5.5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 3											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	10252 ± 0.0	10252	439 ± 0.0	439	221 ± 0.0	221	37.43	37.43 ± 0.0	37.43		
sabre	10373 ± 43.89	10384	1537.19 ± 30.88	1390	155.31 ± 2.59	158	681.17	42.57 ± 0.25	42.4		

name = checkerboard, size = (7.5, 7.5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 0									
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	1211 ± 0.0	1211	30 ± 0.0	30	393 ± 0.0	393	4.76	4.76 ± 0.0	4.76
sabre	4053 ± 23.25	4092	198.25 ± 3.97	187	350.88 ± 3.41	355	14.04	0.88 ± 0.04	0.91

name = checkerboard, size = (7.5, 7.5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 1											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	1211 ± 0.0	1211	30 ± 0.0	30	393 ± 0.0	393	4.49	4.49 ± 0.0	4.49		
sabre	2459 ± 56.57	2107	186.38 ± 11.72	152	284.44 ± 2.91	289	25.29	1.58 ± 0.05	1.64		

name = checkerboard, size = (7.5, 7.5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 2											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	1211 ± 0.0	1211	30 ± 0.0	30	393 ± 0.0	393	4.58	4.58 ± 0.0	4.58		
sabre	2189 ± 48.72	1989	167.81 ± 10.59	138	286.19 ± 3.31	293	57.96	3.62 ± 0.05	3.58		

name = chec	ame = checkerboard, size = (7.5, 7.5), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 3										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	1211 ± 0.0	1211	30 ± 0.0	30	393 ± 0.0	393	4.55	4.55 ± 0.0	4.55		
anhro	1767 + 27 66	1641	125 20 + 6 47	105	270 06 ± 4 72	200	282.04	17 60 + 0 12	17 6		
sabre	1/0/ ± 3/.00	1041	135.30 = 0.47	105	219.06 ± 4.12	290	203.04	17.69 ± 0.12	17.0		

name = checkerboard, size = (7.5, 7.5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 0										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	9289 ± 0.0	9289	219 ± 0.0	219	393 ± 0.0	393	34.14	34.14 ± 0.0	34.14	
sabre	22164 ± 192.21	21331	1027.31 ± 17.21	951	361.69 ± 3.69	377	93.75	5.86 ± 0.21	6.22	

name = checkerboard, size = (7.5, 7.5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 1											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	9289 ± 0.0	9289	219 ± 0.0	219	393 ± 0.0	393	34.49	34.49 ± 0.0	34.49		
sabre	12118 ± 289.99	11694	1276.06 ± 66.48	1103	285.69 ± 3.28	293	118.35	7.4 ± 0.19	7.07		

name = checkerboard, size = (7.5, 7.5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 2										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	9289 ± 0.0	9289	219 ± 0.0	219	393 ± 0.0	393	33.66	33.66 ± 0.0	33.66	
sabre	12194 ± 132.53	11396	1219.62 ± 31.86	1093	287.31 ± 3.31	286	222.24	13.89 ± 0.23	13.07	

name = checkerboard, size = (7.5, 7.5), circuit_type = quantum_simulation, p = 8, repetitions = 16, optimization_level = 3											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	9289 ± 0.0	9289	219 ± 0.0	219	393 ± 0.0	393	33.34	33.34 ± 0.0	33.34		
sabre	11481 ± 103.84	11276	1151.81 ± 30.17	1052	279.0 ± 3.91	278	974.77	60.92 ± 0.4	61.4		

name = checkerboard, size = (7.5, 7.5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 0										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	18521 ± 0.0	18521	435 ± 0.0	435	393 ± 0.0	393	68.34	68.34 ± 0.0	68.34	
sabre	39936 ± 352.99	38867	1884.31 ± 37.54	1805	359.75 ± 3.12	363	169.73	10.61 ± 0.29	10.02	

name = chec	kerboard, size =	(7.5, 7.5), ci	rcuit_type = qua	ntum_simulatio	on, p = 16, repet	titions = 16, opt	imization_level =	1	
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	18521 ± 0.0	18521	435 ± 0.0	435	393 ± 0.0	393	67.22	67.22 ± 0.0	67.22
	02052 + 241 21	02250	0206 00 + 06 0	0054	006 44 + 2 70	200	000 72	12 00 + 0 40	12.0
sabre	23053 = 341.31	23309	2390.00 ± 90.2	2054	200.44 = 3.70	299	222.13	13.92 ± 0.42	15.2

name = chec	name = checkerboard, size = (7.5, 7.5), circuit_type = quantum_simulation, p = 16, repetitions = 16, optimization_level = 2											
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)			
line-graph	18521 ± 0.0	18521	435 ± 0.0	435	393 ± 0.0	393	68.09	68.09 ± 0.0	68.09			
sabre	23286 ± 127.03	22930	2395.75 ± 51.28	2243	288.56 ± 5.34	280	389.15	24.32 ± 0.38	24.34			
name = chec	kerboard, size =	(7.5, 7.5), ci	rcuit_type = quan	tum_simulation	, p = 16, repeti	tions = 16, opti	mization_level =	3				
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)			
line-graph	18521 ± 0.0	18521	435 ± 0.0	435	393 ± 0.0	393	67.2	67.2 ± 0.0	67.2			
sabre	22170 ± 140.89	23060	2352.75 ± 72.43	2121	281.62 ± 4.06	282	1750.38	109.4 ± 0.67	109.23			

7.0.3 Random circuit, kagome and shuriken, against SABRE

```
import line_graph_routing as lgr # Loading these makes these cells stand-alone
import pickle
settings=[]
for name in ['kagome', 'shuriken']:
   for side in range(1,7,2):
        for p in [side**2*500]:
            for optimization_level in [1]:
                setting={'name':name,
                     'size': (side,side),
                     'circuit_type': 'random',
                     'p': p,
                     'repetitions' : 16,
                     'optimization_level' : optimization_level,
                     'methods' : ['sabre']
                    }
                settings.append(setting)
## Uncomment to rerun benchmarks.
#results=[]
#for setting in settings:
   result=lgr.benchmark(**setting)
#
#
    results.append(result)
#
     lgr.print_benchmark(result)
#
#with open('benchmark_data/random.pkl', 'wb') as f:
#
    pickle.dump(results,f)
#Load previously obtained results from disk and show them.
import pickle
with open('benchmark_data/random.pkl','rb') as f:
```

results=pickle.load(f)

for result in results: lgr.print_benchmark(result)

name = kago	ome, size = (1,	1), circuit_typ	e = random, p = :	500, repetit	ions = 16, optimi	zation_level =	1		
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	200 ± 0.0	200	249 ± 0.0	249	12 ± 0.0	12	0.37	0.37 ± 0.0	0.37
sabre	96 ± 0.62	96	207.5 ± 1.19	205	12.0 ± nan	12	1.17	0.07 ± 0.03	0.06
name = kago	ome, size = (3,	3), circuit_typ	be = random, p = .	4500, repeti	tions = 16, optim	ization_level =	1		
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	2611 ± 0.0	2611	568 ± 0.0	568	68 ± 0.0	68	1.14	1.14 ± 0.0	1.14
sabre	1444 ± 25.94	1450	838.88 ± 35.0	720	44.31 ± 1.75	54	10.24	0.64 ± 0.03	0.66
name = kago	ome, size = (5,	5), circuit_typ	e = random, p =	12500, repet	itions = 16, opti	mization_level	- 1		
method	av. n_swaps	min. n_swap	av. depth	min. dep	oth av. n_qubits	min. qubit:	s total time ((s) av. time (s)	min. time (s
line-graph	7592 ± 0.0	7592	735 ± 0.0	7	'35 164 ± 0.0	16	4 3.	.02 3.02 ± 0.0	3.0
sabre	5637 ± 82.31	5452	1774.44 ± 54.79	15	590 108.69 ± 3.56	10	7 32.	72 2.04 ± 0.04	2.0
name = shur	riken, size = (1	, 1), circuit_t	ype = random, p	= 500, repet	titions = 16, opti	mization_level	= 1		
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	122 ± 0.0	122	247 ± 0.0	247	8 ± 0.0	8	0.07	0.07 ± 0.0	0.07
sabre	85 ± nan	85	231.12 ± 0.31	231	8.0 ± nan	8	0.84	1 0.05 ± 0.0	0.05
name = shur	riken, size = (3	, 3), circuit_t	ype = random, p	= 4500, repe	vitions = 16, opt	imization_level	= 1		
method	av. n_swaps	min. n_swap	av. depth	min. dept	h av. n_qubits	min. qubits	total time (s	s) av. time (s)	min. time (s)
line-graph	2262 ± 0.0	2262	415 ± 0.0	41	5 84 ± 0.0	84	1.0	05 1.05 ± 0.0	1.05
sabre	1662 ± 39.26	1447	667.12 ± 39.54	51	.7 65.25 ± 1.38	69	10.5	53 0.66 ± 0.04	0.7

name = shuriken, size = (5, 5), circuit_type = random, p = 12500, repetitions = 16, optimization_level = 1										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	6828 ± 0.0	6828	405 ± 0.0	405	240 ± 0.0	240	2.83	2.83 ± 0.0	2.83	
sabre	6811 ± 192.15	5969	1403.31 ± 94.74	1094	177.31 ± 2.94	176	38.36	2.4 ± 0.04	2.38	

7.0.4 Random circuit, complete graph, against SABRE

```
import line_graph_routing as lgr
import pickle
settings=[]
for name in ['complete']:
   for side in range(3,10,2):
        for p in [side*100]:
            for optimization_level in range(2):
                setting={'name':name,
                         'size': side,
                         'circuit_type': 'random',
                         'p': p,
                         'repetitions' : 16,
                         'optimization_level' : optimization_level,
                         'methods' : ['sabre']
                    }
                settings.append(setting)
## Uncomment to rerun benchmarks
#results=[]
#for setting in settings:
   result=lgr.benchmark(**setting)
#
    results.append(result)
#
    lqr.print_benchmark(result)
#
#with open('benchmark_data/complete.pkl', 'wb') as f:
    pickle.dump(results,f)
#
with open('benchmark_data/complete.pkl','rb') as f:
   results=pickle.load(f)
for result in results:
   lgr.print_benchmark(result)
```


name = complete, size = 3, circuit_type = random, p = 300, repetitions = 16, optimization_level = 0										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	154 ± 0.0	154	334 ± 0.0	334	4 ± 0.0	4	0.12	0.12 ± 0.0	0.12	
sabre	40 ± 0.97	37	267.88 ± 2.34	257	3.0 ± nan	3	1.01	0.06 ± 0.06	0.04	

name = complete, size = 3, circuit_type = random, p = 300, repetitions = 16, optimization_level = 1										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)	
line-graph	140 ± 0.0	140	313 ± 0.0	313	4 ± 0.0	4	0.11	0.11 ± 0.0	0.11	
sabre	34 ± nan	34	220.0 ± nan	220	3.0 ± nan	3	0.97	0.06 ± 0.0	0.07	

name = comp	lete, size = 5,	circuit_type =	random, p = 5	00, repetition	us = 16, optimiza	ation_level = 0			
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	314 ± 0.0	314	548 ± 0.0	548	6 ± 0.0	6	0.19	0.19 ± 0.0	0.19
sabre	88 ± 1.25	84	458.06 ± 4.0	444	5.0 ± nan	5	0.96	0.06 ± nan	0.06

name = comp	name = complete, size = 5, circuit_type = random, p = 500, repetitions = 16, optimization_level = 1										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	286 ± 0.0	286	503 ± 0.0	503	6 ± 0.0	6	0.18	0.18 ± 0.0	0.18		
sabre	72 ± nan	72	387.25 ± 0.31	386	5.0 ± nan	5	1.44	0.09 ± nan	0.09		

name = comp	lete, size = 7,	circuit_type =	random, p = 70	0, repetitions	a = 16, optimiza	tion_level = 0			
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	436 ± 0.0	436	723 ± 0.0	723	8 ± 0.0	8	0.25	0.25 ± 0.0	0.25
sabre	115 ± 1.59	115	609.25 ± 6.01	586	7.0 ± nan	7	1.65	0.1 ± 0.02	0.09

name = comp	lete, size = 7,	circuit_type =	random, $p = 70$	0, repetitions	= 16, optimizat	tion_level = 1			
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
	470 1 0 0	470	707	207			0.00		0.00
line-graph	478 ± 0.0	478	787 ± 0.0	/8/	8 ± 0.0	8	0.26	0.26 ± 0.0	0.26
anhra	110 + 0 20	117	572 OF ± 1 66	567	7.0 ± non	7	2.24	0 14 + 0 01	0.12
Sable	110 - 0.30	117	575.25 - 1.00	507	7.0 ± IIali	'	2.24	0.14 - 0.01	0.13

name = comp	olete, size = 9,	circuit_type =	random, p = 90	00, repetition	s = 16, optimiza	tion_level = 0			
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	593 ± 0.0	593	946 ± 0.0	946	10 ± 0.0	10	0.43	0.43 ± 0.0	0.43
sabre	151 ± 2.0	143	770.0 ± 6.84	730	9.0 ± nan	9	1.6	0.1 ± nan	0.1
name = comp	olete, size = 9,	circuit_type =	random, p = 90	00, repetition	s = 16, optimiza	tion_level = 1			
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	650 ± 0.0	650	1036 ± 0.0	1036	10 ± 0.0	10	0.33	0.33 ± 0.0	0.33
sabre	149 ± 0.59	146	726.5 ± 1.59	720	9.0 ± nan	9	2.74	0.17 ± 0.02	0.16

7.0.5 Against other routing methods

Above, we only ran SABRE because it outperforms the other methods available in Qiskit by default. The standard methods available are

```
from qiskit import transpiler
transpiler.preset_passmanagers.plugin.list_stage_plugins('routing')
```

```
['basic', 'lookahead', 'none', 'sabre', 'stochastic']
```

```
import line_graph_routing as lgr
import pickle
settings=[]
for name in ['kagome', 'shuriken']:
    for side in [1]:
        for p in [1]:
            for optimization_level in [3]:
                setting={'name':name,
                         'size': (side,side),
                         'circuit_type': 'quantum_simulation',
                         'p': p,
                         'repetitions' : 16,
                         'optimization_level' : optimization_level,
                         'methods' : ['basic', 'lookahead', 'sabre', 'stochastic']
                    }
                settings.append(setting)
## Uncomment to rerun benchmarks
#results=[]
#for setting in settings:
#
   result=lqr.benchmark(**setting)
#
  results.append(result)
     lqr.print_benchmark(result)
#
```

```
#with open('benchmark_data/other_methods_1x1.pkl','wb') as f:
# pickle.dump(results,f)
with open('benchmark_data/other_methods_1x1.pkl','rb') as f:
    results=pickle.load(f)
```

for result in results:
 lgr.print_benchmark(result)

name = kago	aame = kagome, size = (1, 1), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 3										
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)		
line-graph	12 ± 0.0	12	7 ± 0.0	7	12 ± 0.0	12	0.08	0.08 ± 0.0	0.08		
basic	14 ± 0.0	14	17.0 ± 0.0	17	8.0 ± 0.0	8	0.05	0.05 ± 0.0	0.05		
lookahead	8 ± nan	8	8.81 ± 0.19	8	8.0 ± nan	8	19.55	1.22 ± 0.04	1.26		
sabre	6 ± nan	6	6.0 ± nan	6	8.0 ± nan	8	0.49	0.03 ± 0.0	0.03		
stochastic	6 ± 0.62	6	6.06 ± 0.16	6	8.0 ± nan	8	0.62	0.04 ± 0.0	0.03		

now a logare size = (1, 1) simulations a substant simulation n = 1 constitutions = 16 entitication logal = 2

name = shur	name = shuriken, size = (1, 1), circuit_type = quantum_simulation, p = 1, repetitions = 16, optimization_level = 3											
where α is some since one of doubt wire doubt on public since α by the set of the (α) and the (α)												
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)			
line-graph	8 ± 0.0	8	9 ± 0.0	9	8 ± 0.0	8	0.05	0.05 ± 0.0	0.05			
basic	16 ± 0.0	16	19.0 ± 0.0	19	8.0 ± 0.0	8	0.04	0.04 ± 0.0	0.04			
lookahead	6 ± nan	6	8.0 ± nan	8	8.0 ± nan	8	14.22	0.89 ± 0.01	0.88			
sabre	6 ± nan	6	7.62 ± 0.22	7	8.0 ± nan	8	0.5	0.03 ± 0.0	0.04			
stochastic	6 ± nan	6	7.0 ± nan	7	8.0 ± nan	8	0.58	0.04 ± 0.0	0.04			

The method 'lookahead' takes an impractical amount of time, so we exclude it when running benchmarks for larger/deeper circuits.

```
import line_graph_routing as lgr
import pickle
settings=[]
for name in ['kagome','shuriken']:
    for side in [3]:
        for p in [3]:
            for optimization_level in [3]:
               setting={'name':name,
                   'size': (side,side),
                   'circuit_type': 'quantum_simulation',
                   'p': p,
                   'repetitions' : 16,
                   'optimization_level' : optimization_level,
```

```
'methods' : ['basic', 'sabre', 'stochastic']
                    }
                settings.append(setting)
## Uncomment to rerun benchmarks
#results=[]
#for setting in settings:
     result=lqr.benchmark(**setting)
#
#
    results.append(result)
#
     lqr.print_benchmark(result)
#
#with open('benchmark_data/other_methods_3x3.pkl', 'wb') as f:
  pickle.dump(results,f)
#
with open('benchmark_data/other_methods_3x3.pkl','rb') as f:
    results=pickle.load(f)
for result in results:
    lgr.print_benchmark(result)
```

name = kago	me, size = (3, 3), circuit_typ	e = quantum_si	mulation, p =	3, repetitions	= 16, optimizatio	n_level = 3		
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	336 ± 0.0	336	43 ± 0.0	43	68 ± 0.0	68	1.33	1.33 ± 0.0	1.33
basic	475 ± 0.0	475	287.0 ± 0.0	287	45.0 ± 0.0	45	1.21	1.21 ± 0.0	1.21
sabre	215 ± 3.56	203	63.25 ± 3.78	50	43.94 ± 0.97	45	12.37	0.77 ± 0.04	0.93

71 46.44 ± 1.5

name = shur	iken, size = (3,	3), circuit_t	ype = quantum_	simulation, p	= 3, repetitions	s = 16, optimizat	tion_level = 3		
method	av. n_swaps	min. n_swap	av. depth	min. depth	av. n_qubits	min. qubits	total time (s)	av. time (s)	min. time (s)
line-graph	390 ± 0.0	390	34 ± 0.0	34	84 ± 0.0	84	1.81	1.81 ± 0.0	1.81
basic	1195 ± 0.0	1195	491.0 ± 0.0	491	71.0 ± 0.0	71	2.14	2.14 ± 0.0	2.14
sabre	380 ± 10.03	344	78.19 ± 4.42	66	64.19 ± 0.94	67	24.94	1.56 ± 0.12	2.31
stochastic	813 ± 32.04	609	104.0 ± 7.88	76	66.56 ± 1.44	71	62.25	3.89 ± 0.09	3.76

We see SABRE outperforms the other methods available by default in Qiskit, but not line-graph routing for the circuits considered.

48

28.09 1.76 ± 0.07

1.46

7.0.6 Wall-clock time of line-graph routing

428 85.88 ± 4.12

stochastic 410 ± 14.31

Create a random graph, construct the line graph, create a circuit on the line graph, and put this circuit into line-graph routing

```
from time import time
side=25
lg = lgr.kagome(side, side)
print('number of nodes =',lg.number_of_nodes())
```

```
La=10**5
print('number of gates =',La)
qc = lgr.random_circuit(lg, La)
begin=time()
qc = lgr.line_graph_route(qc)
end=time()
print('wall clock time =',end-begin,'(s)')
```

number of nodes = 1976
number of gates = 100000
wall clock time = 25.544434070587158 (s)